

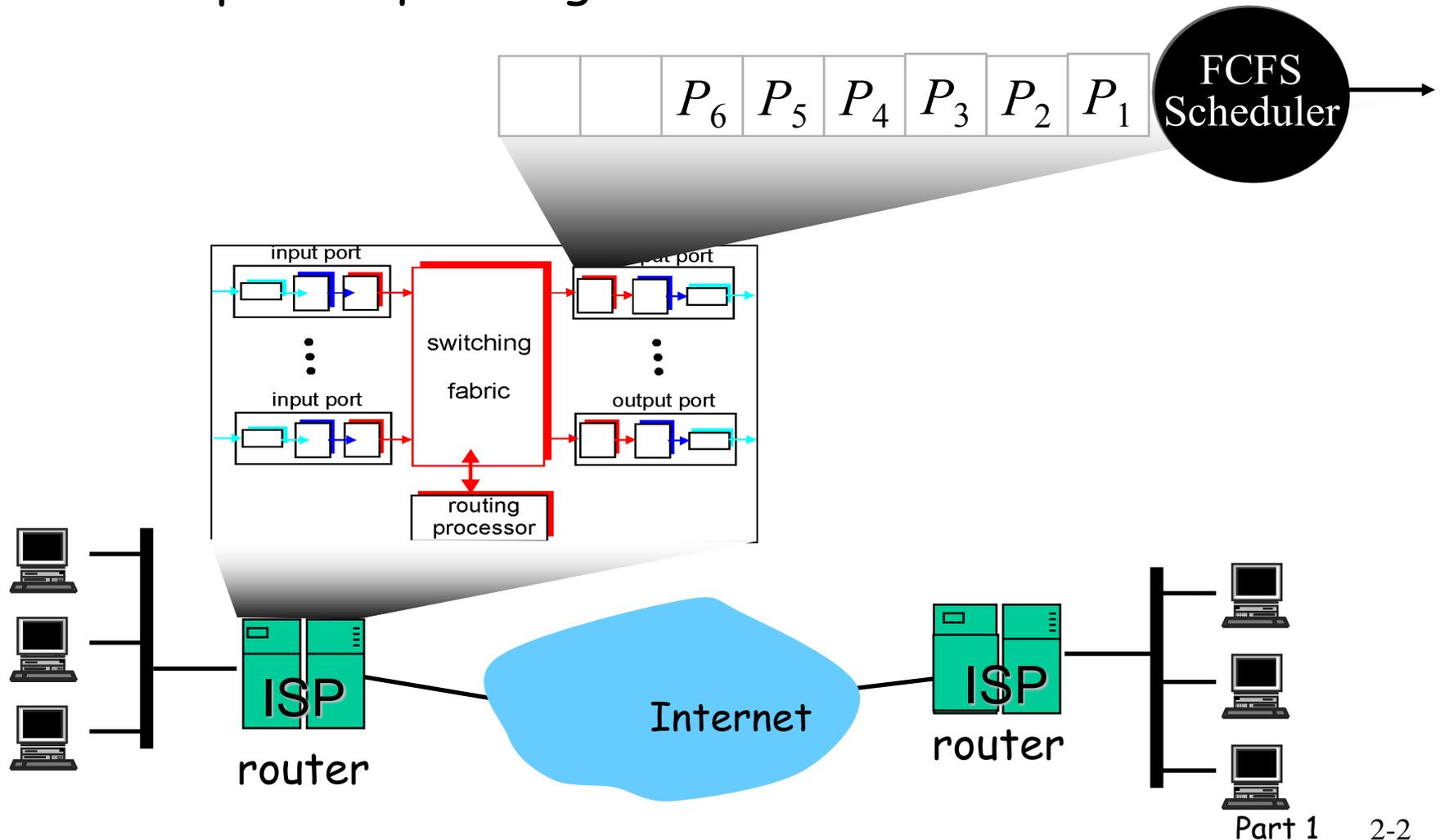
Randomization in Router Queue Management

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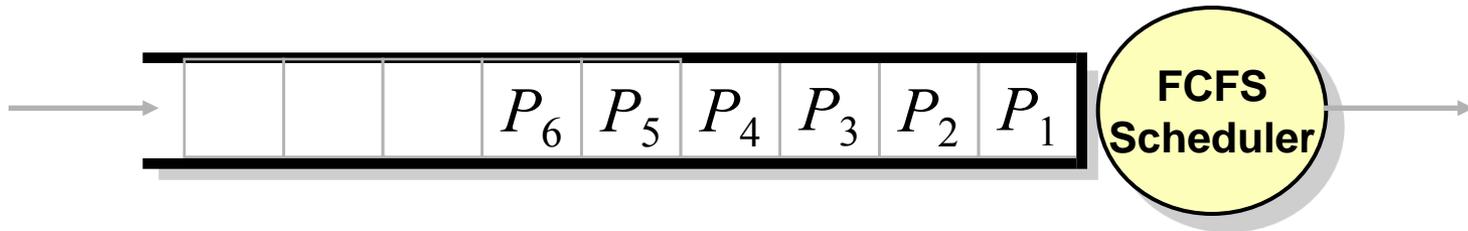
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Randomization in Router Queue Management

- normally, packets dropped only when queue overflows
 - "Drop-tail" queueing

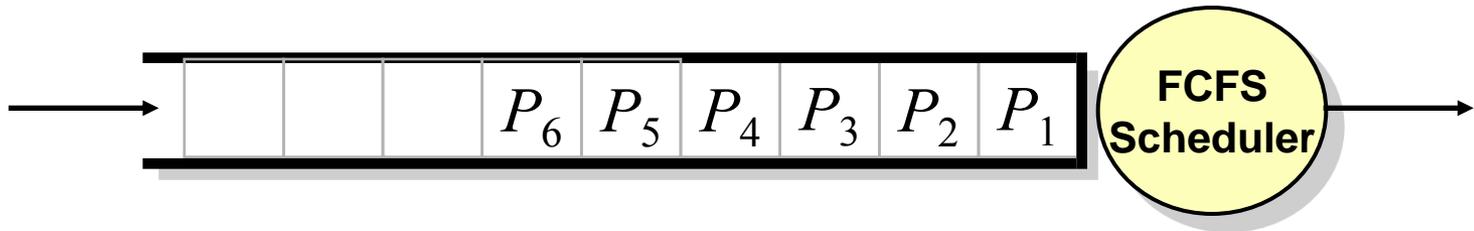


The case against drop-tail queue management



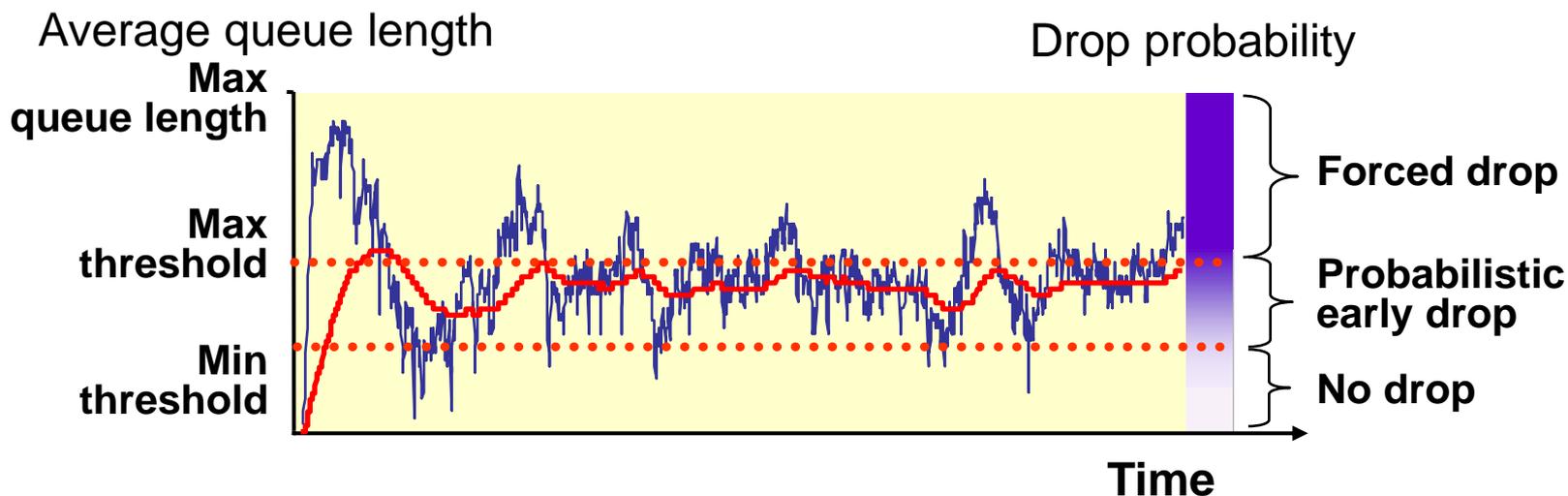
- ❑ large queues in routers are “a bad thing”
 - End-to-end latency dominated by length of queues at switches in network
- ❑ allowing queues to overflow is “a bad thing”
 - connections transmitting at high rates can starve connections transmitting at low rates
 - connections can *synchronize* their response to congestion

Idea: early random packet drop



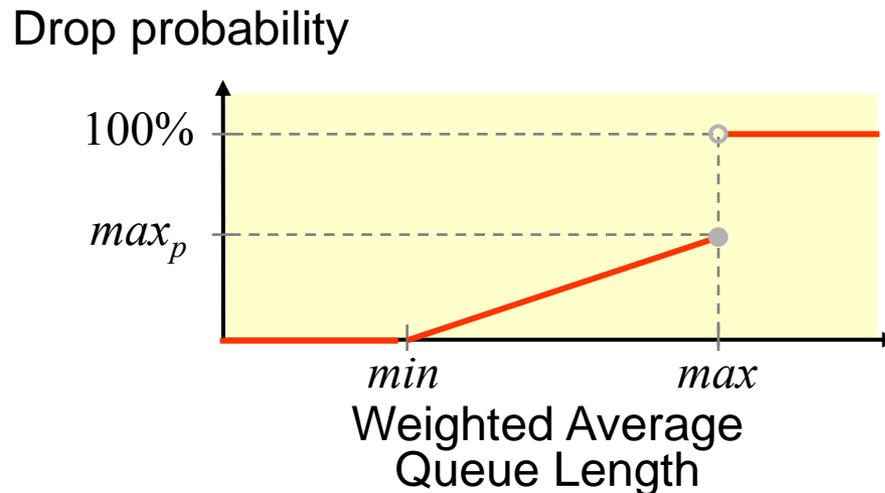
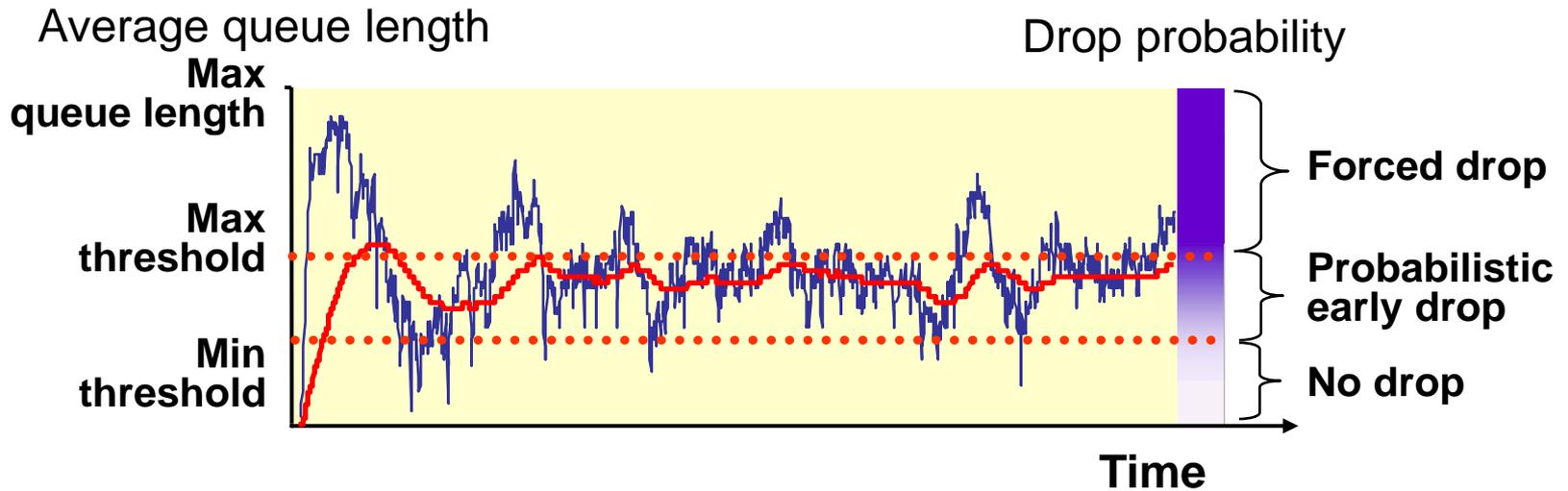
- When queue length exceeds threshold, packets dropped with fixed *probability*
 - probabilistic packet drop: flows see same loss *rate*
 - problem: bursty traffic (burst arrives when queue is near full) can be overpenalized

Random early detection (RED) packet drop



- use *exponential average* of queue length to determine when to drop
 - avoid overly penalizing short-term bursts
 - React to longer term trends
- tie drop prob. to weighted avg. queue length
 - avoids over-reaction to mild overload conditions

Random early detection (RED) packet drop



Random early detection (RED) packet drop

- ❑ large number (5) of parameters: difficult to tune (at least for http traffic)
- ❑ gains over drop-tail FCFS not that significant
- ❑ still not widely deployed ...

RED: why probabilistic drop?

- ❑ provide gentle transition from no-drop to all-drop
 - provide “gentle” early warning
- ❑ provide same loss rate to all sessions:
 - with tail-drop, low-sending-rate sessions can be completely starved
- ❑ avoid synchronized loss bursts among sources
 - avoid cycles of large-loss followed by no-transmission