

Chapter 4

Network Layer

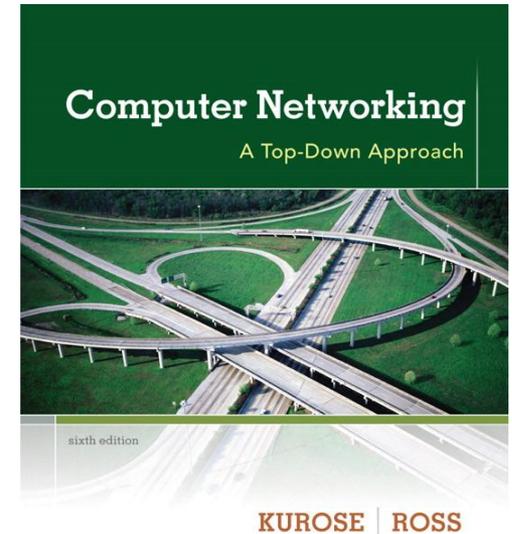
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*Computer
Networking: A
Top Down
Approach*
6th edition
Jim Kurose, Keith Ross
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Chapter 4: outline

4.1 introduction

4.2 virtual circuit and datagram networks

4.3 what's inside a router

4.4 IP: Internet Protocol

- datagram format
- IPv4 addressing
- ICMP
- IPv6

4.5 routing algorithms

- link state
- distance vector
- hierarchical routing

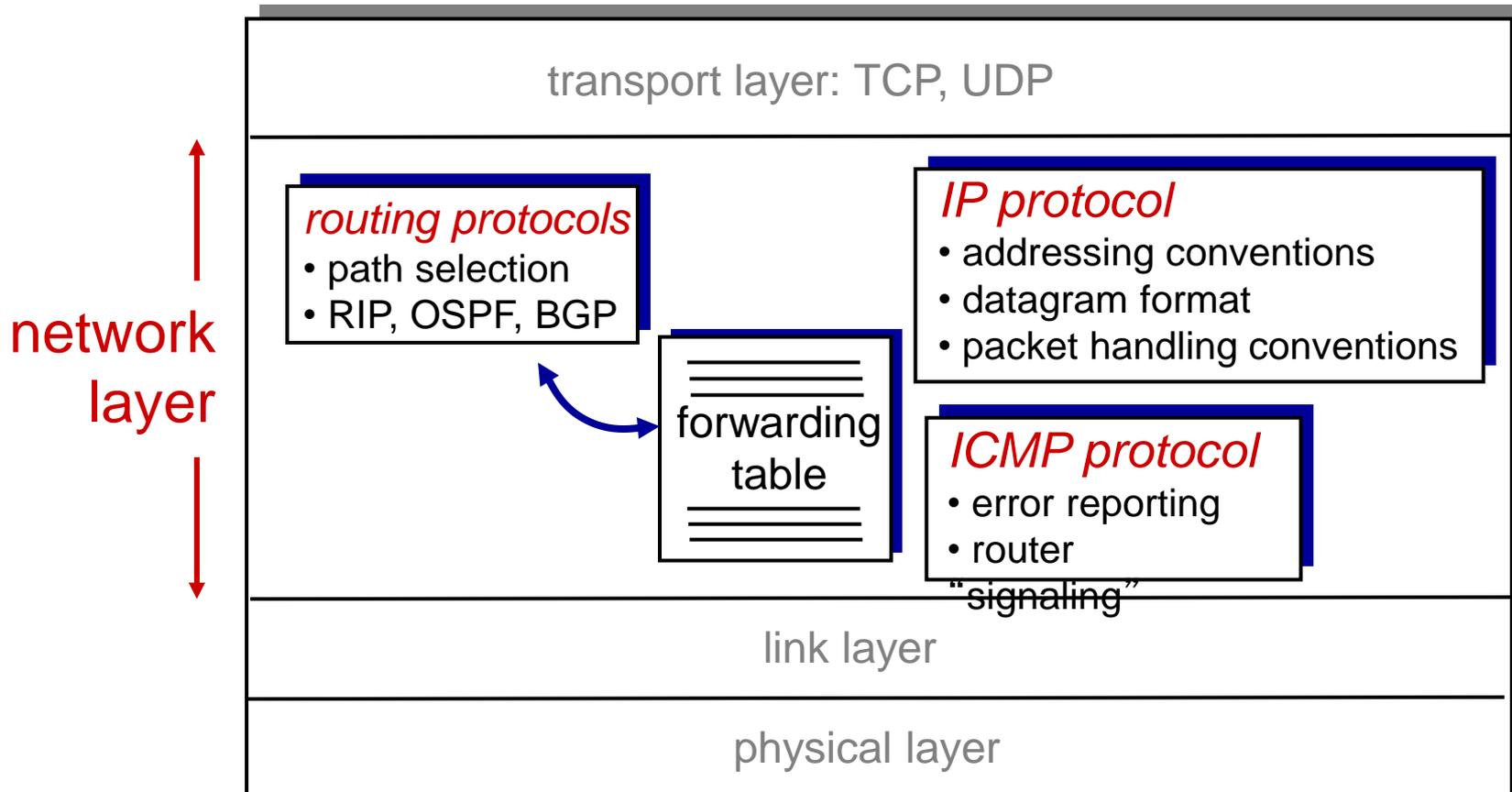
4.6 routing in the Internet

- RIP
- OSPF
- BGP

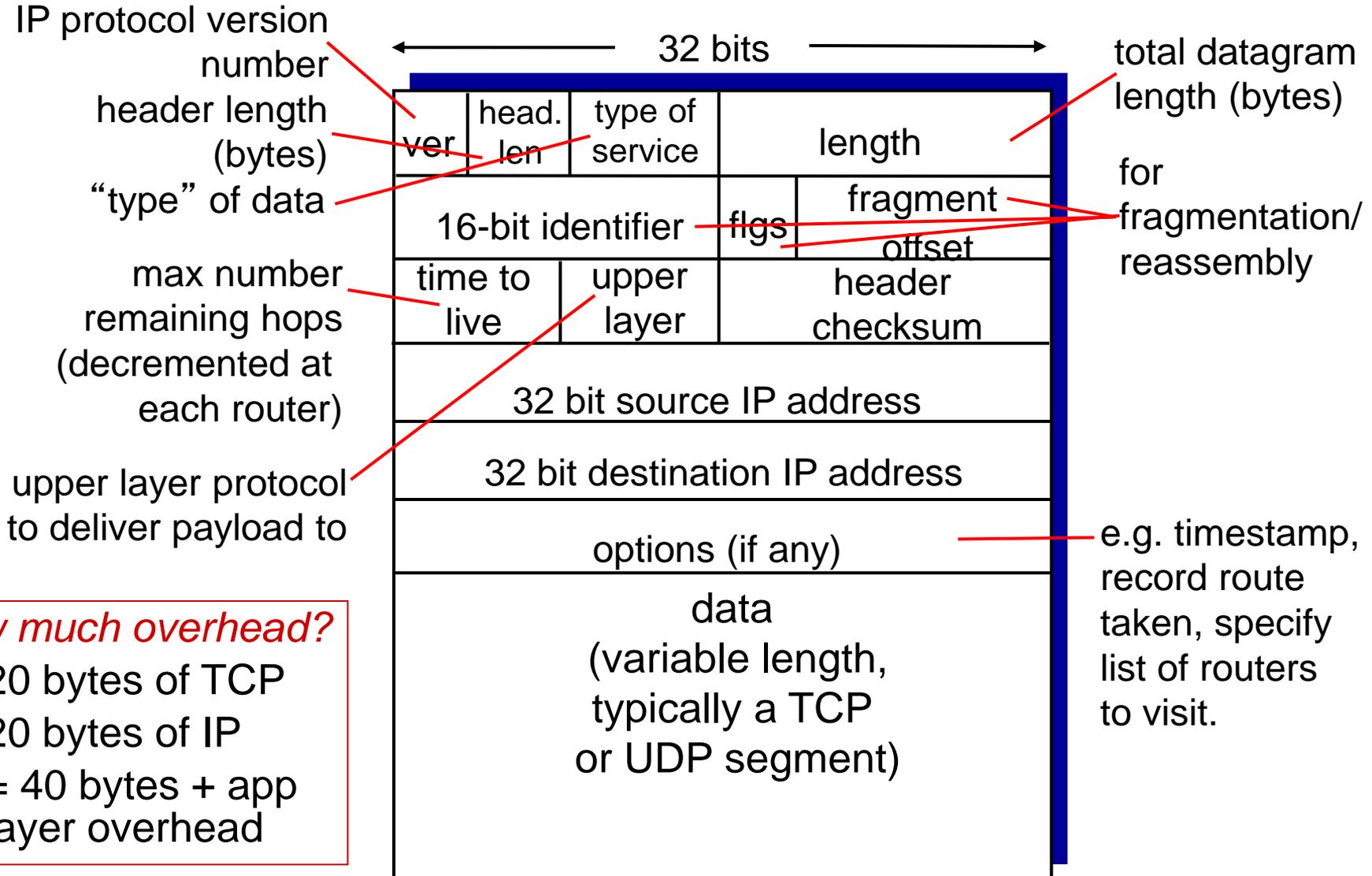
4.7 broadcast and multicast routing

The Internet network layer

host, router network layer functions:



IP datagram format

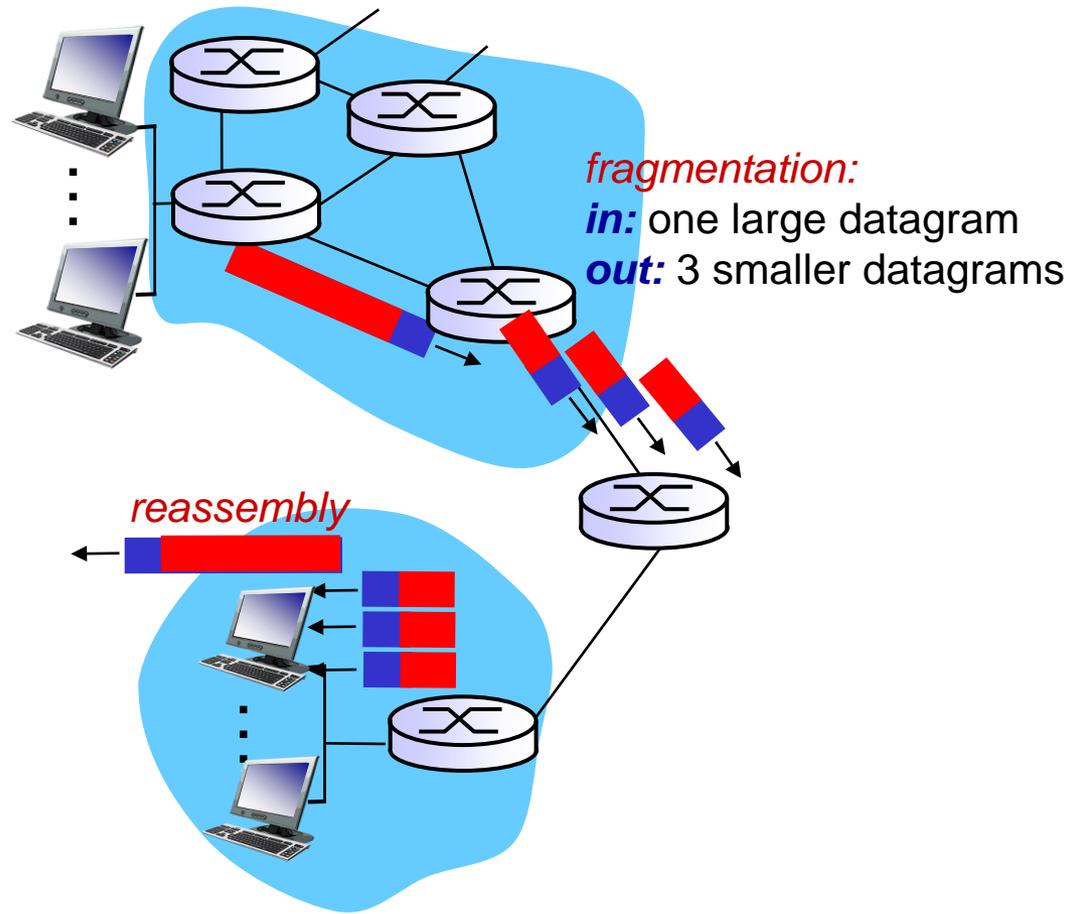


how much overhead?

- ❖ 20 bytes of TCP
- ❖ 20 bytes of IP
- ❖ = 40 bytes + app layer overhead

IP fragmentation, reassembly

- ❖ network links have MTU (max.transfer size) - largest possible link-level frame
 - different link types, different MTUs
- ❖ large IP datagram divided (“fragmented”) within net
 - one datagram becomes several datagrams
 - “reassembled” only at final destination
 - IP header bits used to identify, order related fragments



IP fragmentation, reassembly

example:

- ❖ 4000 byte datagram
- ❖ MTU = 1500 bytes

	length	ID	fragflag	offset	
	=4000	=x	=0	=0	

one large datagram becomes several smaller datagrams

1480 bytes in
data field

offset =
1480/8

	length	ID	fragflag	offset	
	=1500	=x	=1	=0	

	length	ID	fragflag	offset	
	=1500	=x	=1	=185	

	length	ID	fragflag	offset	
	=1040	=x	=0	=370	

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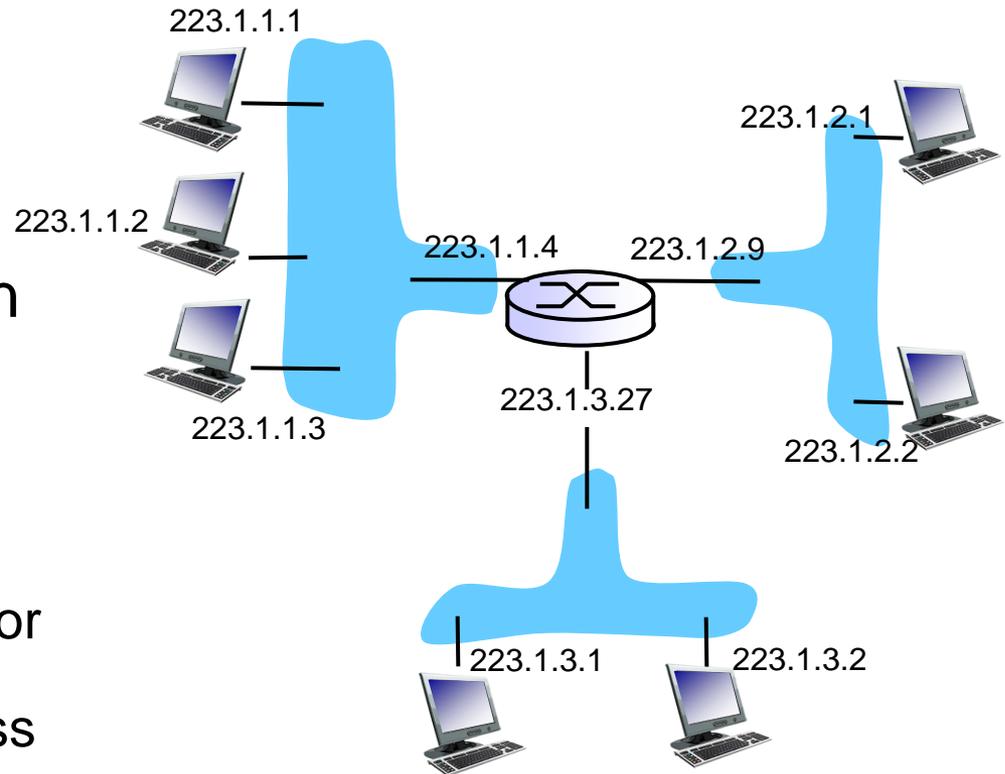
4.6 routing in the Internet

- RIP
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IP addressing: introduction

- ❖ **IP address:** 32-bit identifier for host, router *interface*
- ❖ **interface:** connection between host/router and physical link
 - router's typically have multiple interfaces
 - host typically has one or two interfaces (e.g., wired Ethernet, wireless 802.11)
- ❖ **IP addresses associated with each interface**



$$223.1.1.1 = \underbrace{11011111}_{223} \underbrace{00000001}_{1} \underbrace{00000001}_{1} \underbrace{00000001}_{1}$$

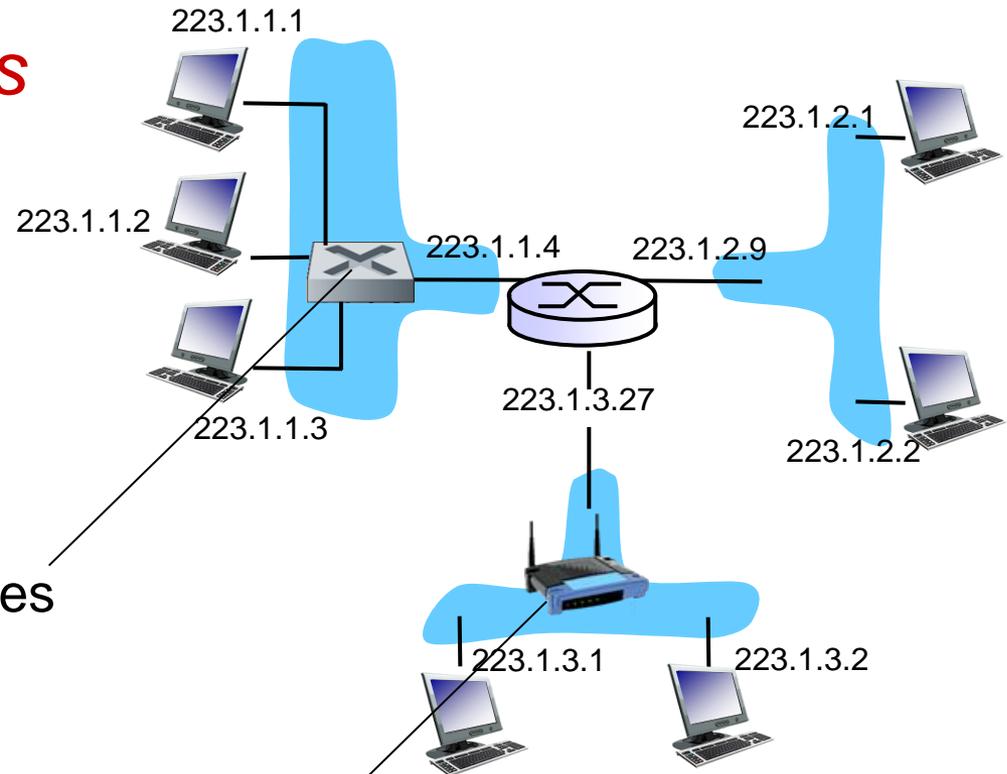
IP addressing: introduction

Q: *how are interfaces actually connected?*

A: *we'll learn about that in chapter 5, 6.*

A: wired Ethernet interfaces connected by Ethernet switches

For now: don't need to worry about how one interface is connected to another (with no intervening router)



A: wireless WiFi interfaces connected by WiFi base station

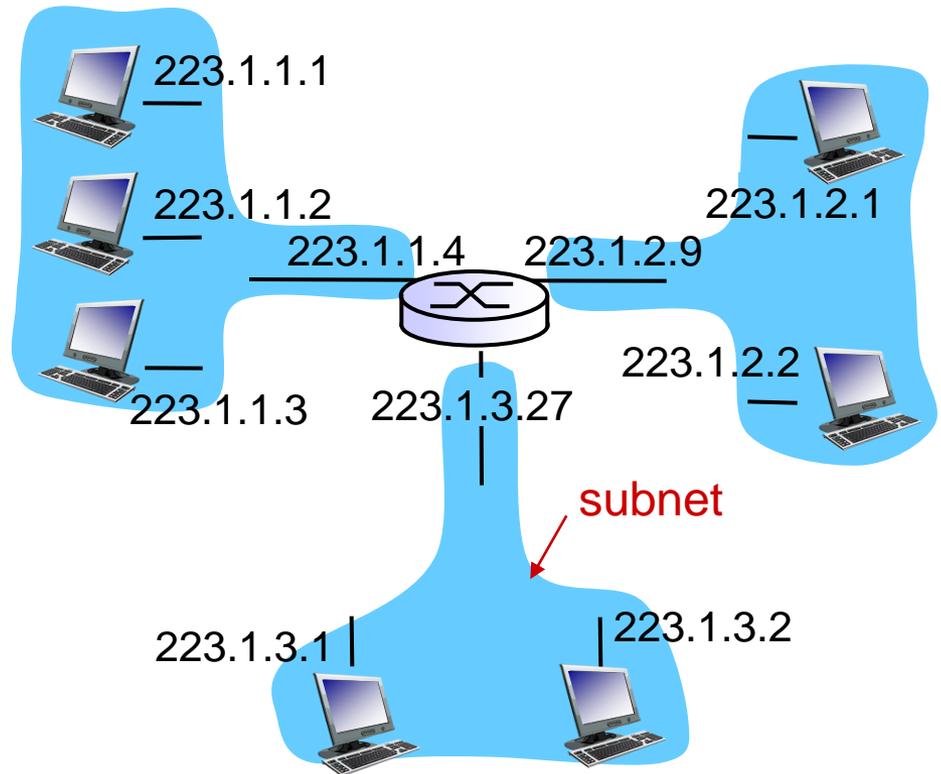
Subnets

❖ IP address:

- subnet part - high order bits
- host part - low order bits

❖ *what's a subnet ?*

- device interfaces with same subnet part of IP address
- can physically reach each other *without intervening router*

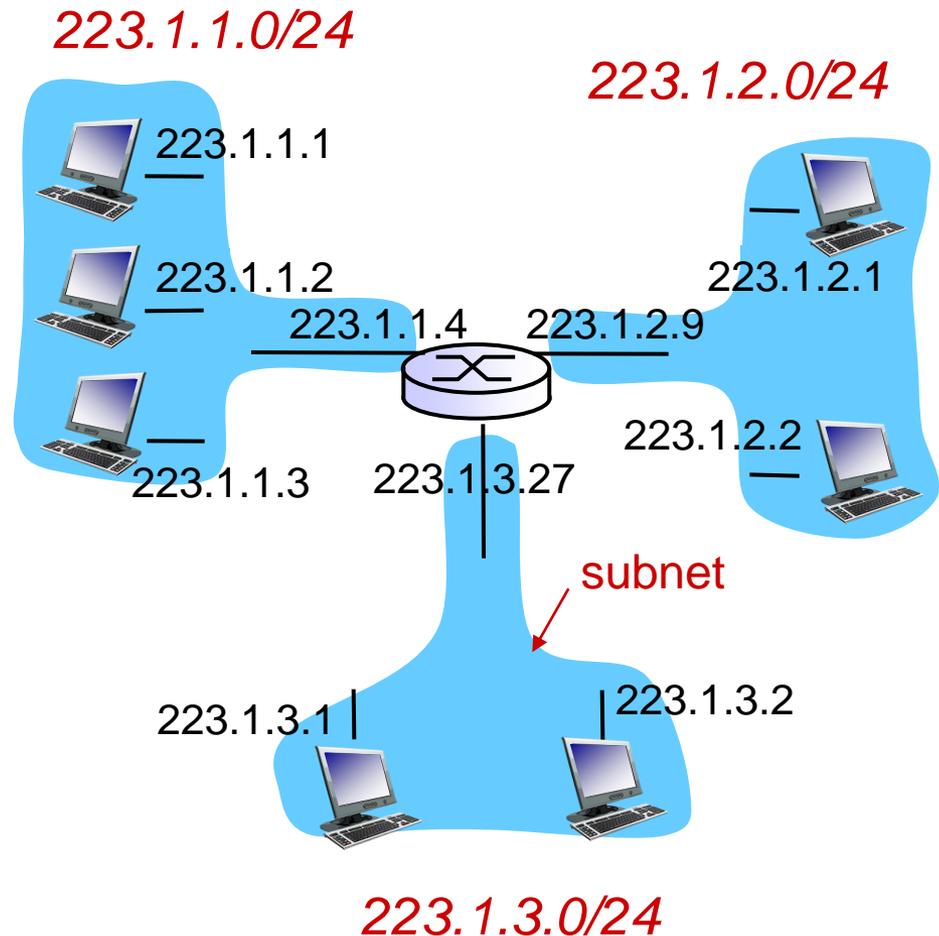


network consisting of 3 subnets

Subnets

recipe

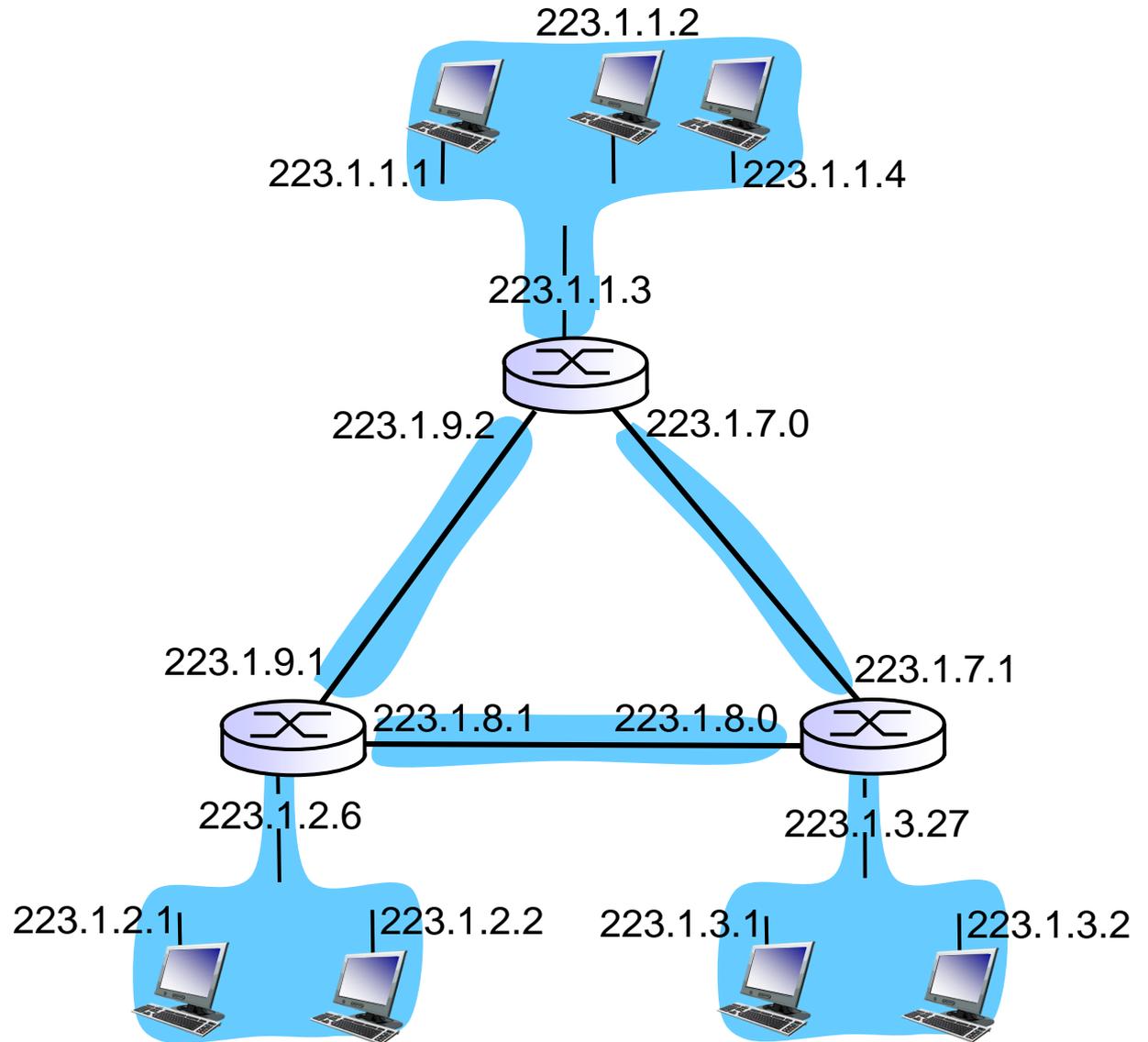
- ❖ to determine the subnets, detach each interface from its host or router, creating islands of isolated networks
- ❖ each isolated network is called a *subnet*



subnet mask: /24

Subnets

how many?



IP addressing: CIDR

CIDR: Classless InterDomain Routing

- subnet portion of address of arbitrary length
- address format: **a.b.c.d/x**, where x is # bits in subnet portion of address



IP addresses: how to get one?

Q: How does a *host* get IP address?

- ❖ hard-coded by system admin in a file
 - Windows: control-panel->network->configuration->tcp/ip->properties
 - UNIX: /etc/rc.config
- ❖ **DHCP: Dynamic Host Configuration Protocol:** dynamically get address from as server
 - “plug-and-play”

DHCP: Dynamic Host Configuration Protocol

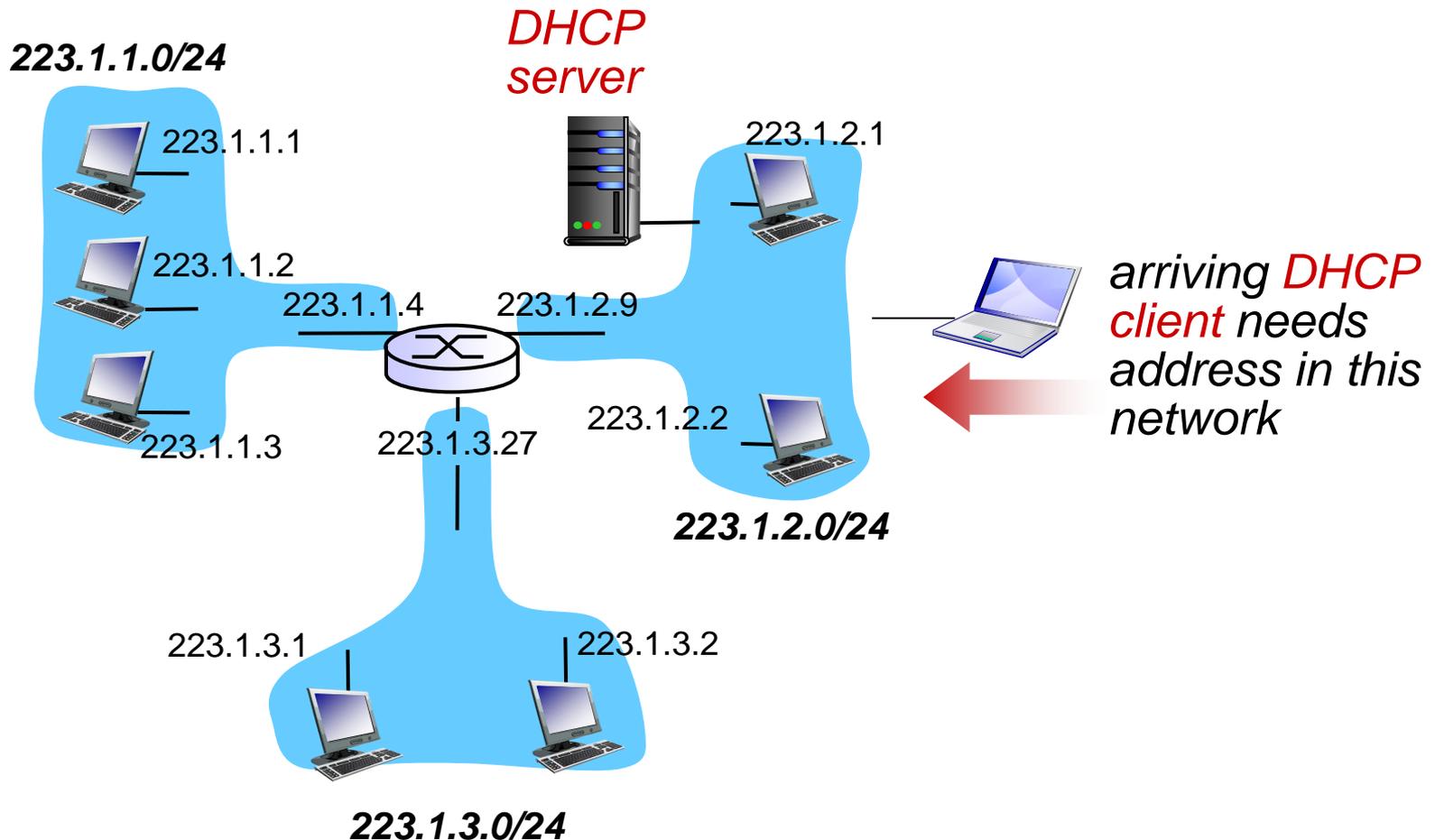
goal: allow host to *dynamically* obtain its IP address from network server when it joins network

- can renew its lease on address in use
- allows reuse of addresses (only hold address while connected/“on”)
- support for mobile users who want to join network (more shortly)

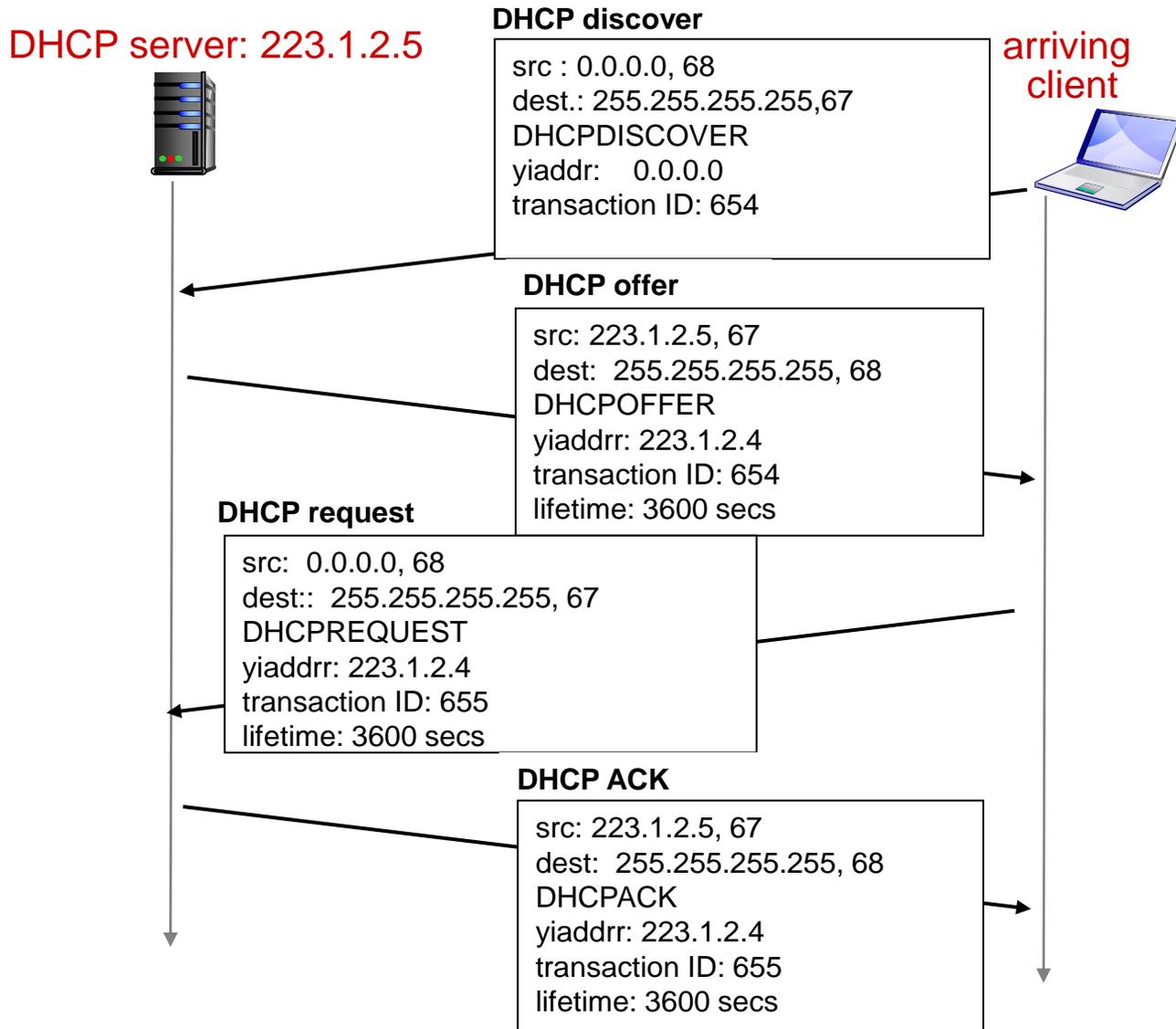
DHCP overview:

- host broadcasts “DHCP discover” msg [optional]
- DHCP server responds with “DHCP offer” msg [optional]
- host requests IP address: “DHCP request” msg
- DHCP server sends address: “DHCP ack” msg

DHCP client-server scenario



DHCP client-server scenario

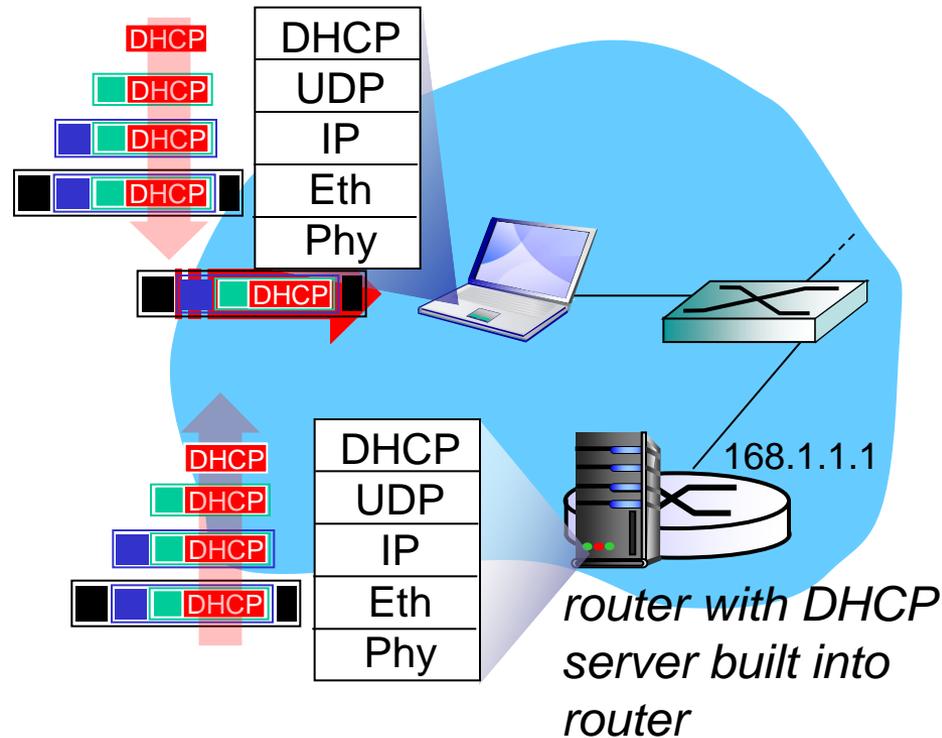


DHCP: more than IP addresses

DHCP can return more than just allocated IP address on subnet:

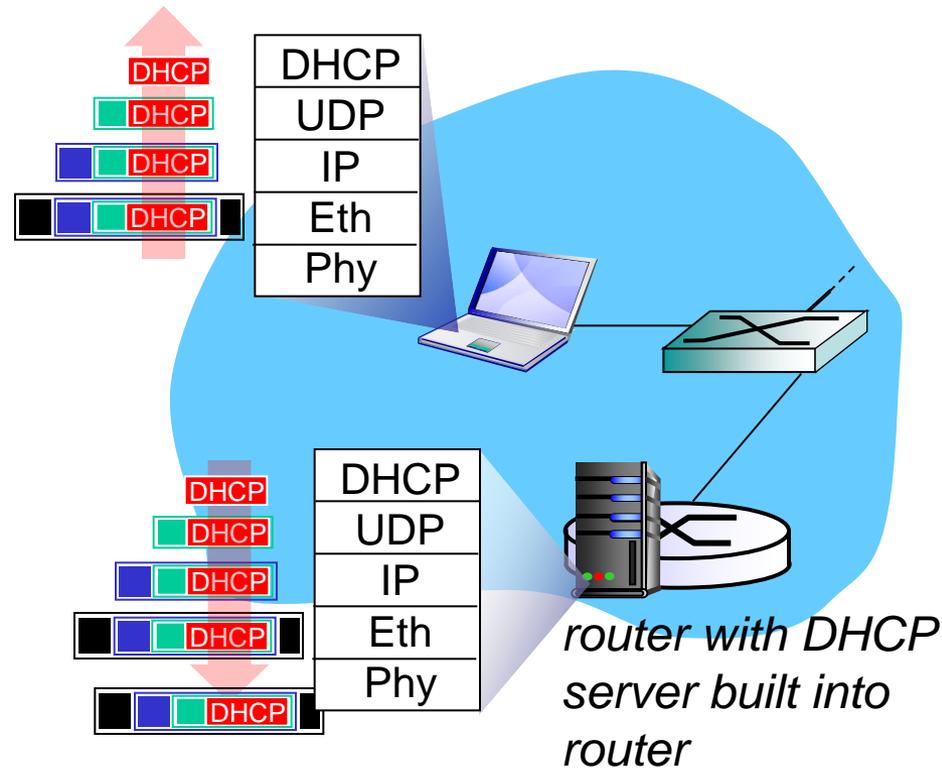
- address of first-hop router for client
- name and IP address of DNS sever
- network mask (indicating network versus host portion of address)

DHCP: example



- ❖ connecting laptop needs its IP address, addr of first-hop router, addr of DNS server: use DHCP
- ❖ DHCP request encapsulated in UDP, encapsulated in IP, encapsulated in 802.1 Ethernet
- ❖ Ethernet frame broadcast (dest: FFFFFFFFFFFFFFFF) on LAN, received at router running DHCP server
- ❖ Ethernet demuxed to IP demuxed, UDP demuxed to DHCP

DHCP: example



- ❖ DHCP server formulates DHCP ACK containing client's IP address, IP address of first-hop router for client, name & IP address of DNS server
- ❖ encapsulation of DHCP server, frame forwarded to client, demuxing up to DHCP at client
- ❖ client now knows its IP address, name and IP address of DNS server, IP address of its first-hop router

DHCP: Wireshark output (home LAN)

Message type: **Boot Request (1)**

Hardware type: Ethernet
Hardware address length: 6
Hops: 0

Transaction ID: 0x6b3a11b7

Seconds elapsed: 0

Bootp flags: 0x0000 (Unicast)

Client IP address: 0.0.0.0 (0.0.0.0)

Your (client) IP address: 0.0.0.0 (0.0.0.0)

Next server IP address: 0.0.0.0 (0.0.0.0)

Relay agent IP address: 0.0.0.0 (0.0.0.0)

Client MAC address: Wistron_23:68:8a (00:16:d3:23:68:8a)

Server host name not given

Boot file name not given

Magic cookie: (OK)

Option: (t=53,l=1) **DHCP Message Type = DHCP Request**

Option: (61) Client identifier

Length: 7; Value: 010016D323688A;

Hardware type: Ethernet

Client MAC address: Wistron_23:68:8a (00:16:d3:23:68:8a)

Option: (t=50,l=4) Requested IP Address = 192.168.1.101

Option: (t=12,l=5) Host Name = "nomad"

Option: (55) Parameter Request List

Length: 11; Value: 010F03062C2E2F1F21F92B

1 = Subnet Mask; 15 = Domain Name

3 = Router; 6 = Domain Name Server

44 = NetBIOS over TCP/IP Name Server

.....

request

Message type: **Boot Reply (2)**

Hardware type: Ethernet

Hardware address length: 6

Hops: 0

Transaction ID: 0x6b3a11b7

Seconds elapsed: 0

Bootp flags: 0x0000 (Unicast)

Client IP address: 192.168.1.101 (192.168.1.101)

Your (client) IP address: 0.0.0.0 (0.0.0.0)

Next server IP address: 192.168.1.1 (192.168.1.1)

Relay agent IP address: 0.0.0.0 (0.0.0.0)

Client MAC address: Wistron_23:68:8a (00:16:d3:23:68:8a)

Server host name not given

Boot file name not given

Magic cookie: (OK)

Option: (t=53,l=1) DHCP Message Type = DHCP ACK

Option: (t=54,l=4) Server Identifier = 192.168.1.1

Option: (t=1,l=4) Subnet Mask = 255.255.255.0

Option: (t=3,l=4) Router = 192.168.1.1

Option: (6) Domain Name Server

Length: 12; Value: 445747E2445749F244574092;

IP Address: 68.87.71.226;

IP Address: 68.87.73.242;

IP Address: 68.87.64.146

Option: (t=15,l=20) Domain Name = "hsd1.ma.comcast.net."

reply

IP addresses: how to get one?

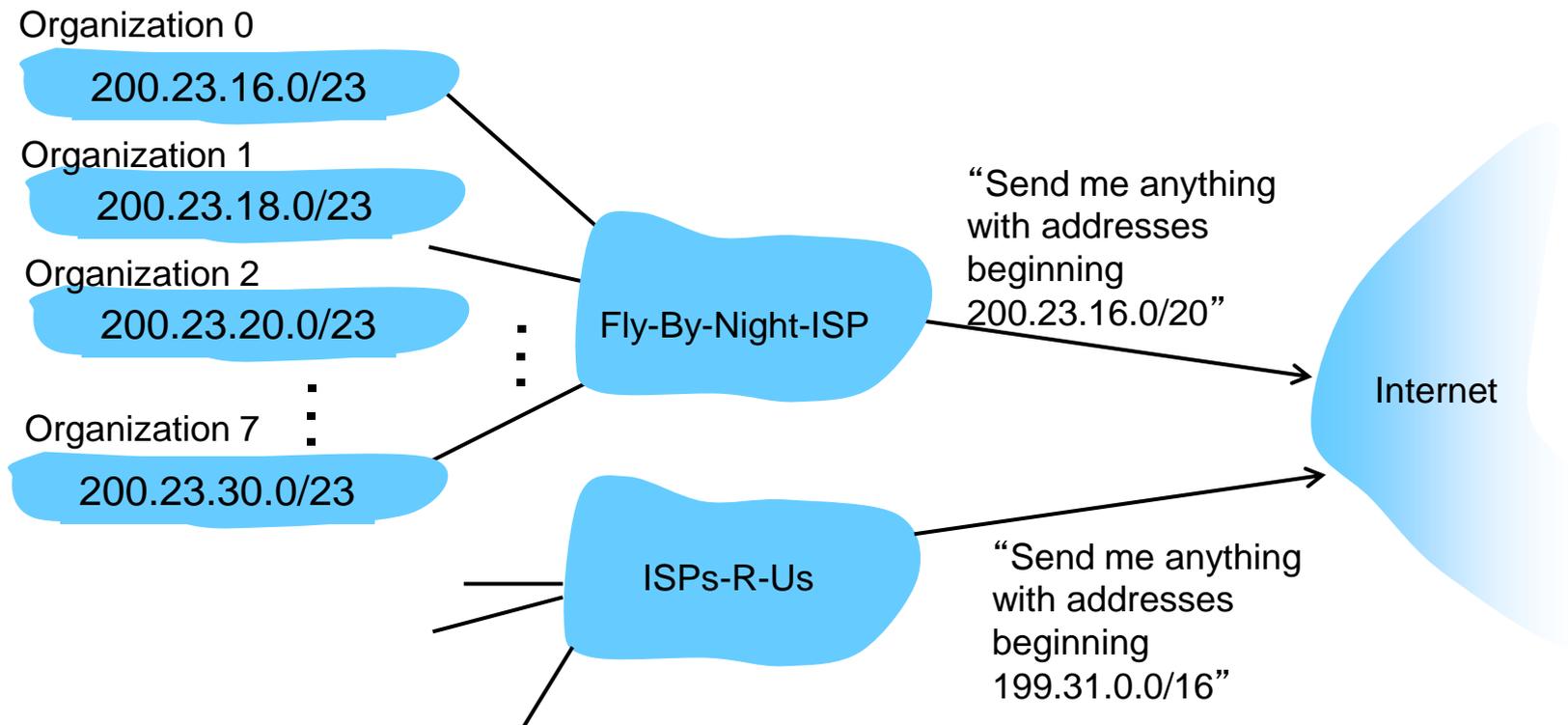
Q: how does *network* get subnet part of IP addr?

A: gets allocated portion of its provider ISP's address space

ISP's block	<u>11001000</u>	<u>00010111</u>	<u>00010000</u>	00000000	200.23.16.0/20
Organization 0	<u>11001000</u>	<u>00010111</u>	<u>00010000</u>	00000000	200.23.16.0/23
Organization 1	<u>11001000</u>	<u>00010111</u>	<u>00010010</u>	00000000	200.23.18.0/23
Organization 2	<u>11001000</u>	<u>00010111</u>	<u>00010100</u>	00000000	200.23.20.0/23
...
Organization 7	<u>11001000</u>	<u>00010111</u>	<u>00011110</u>	00000000	200.23.30.0/23

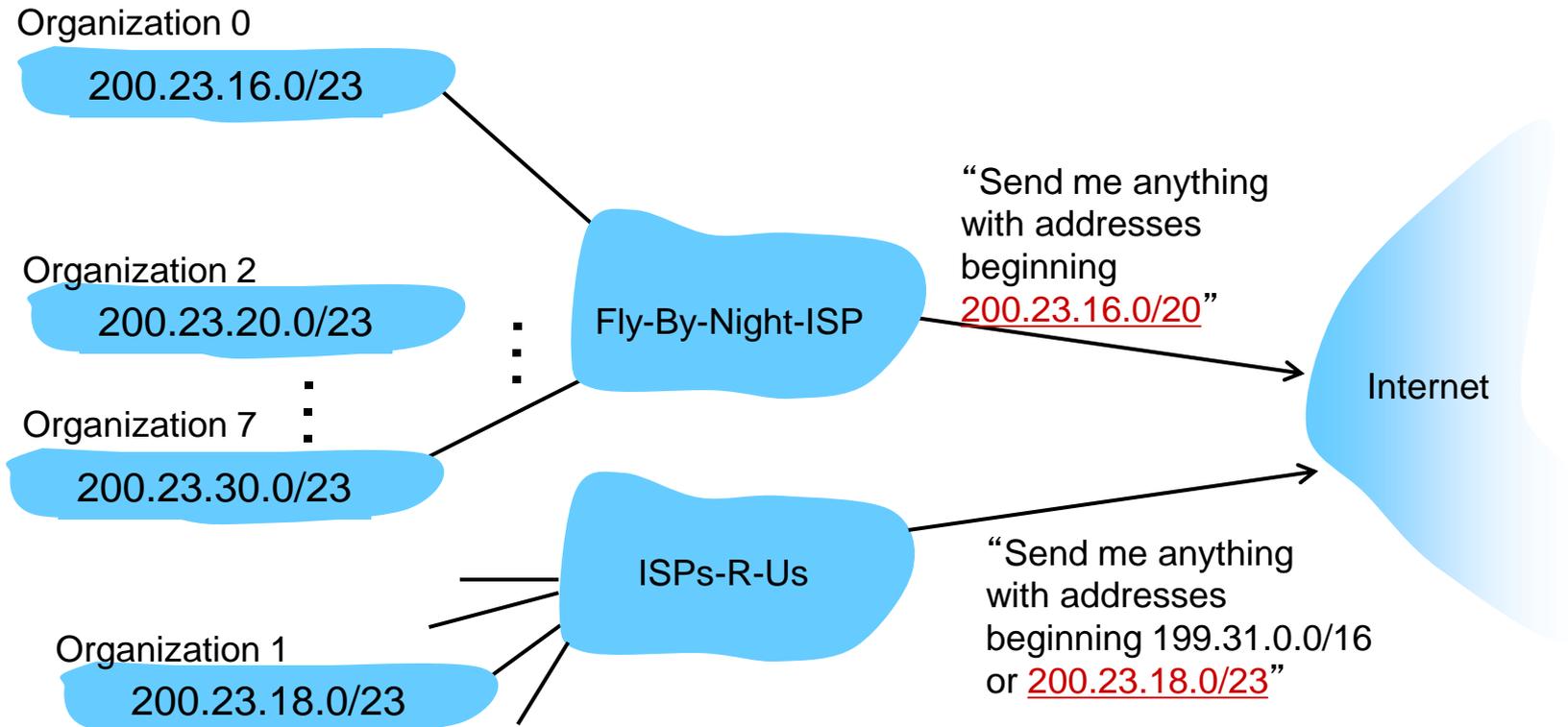
Hierarchical addressing: route aggregation

hierarchical addressing allows efficient advertisement of routing information:



Hierarchical addressing: more specific routes

ISPs-R-U's has a more specific route to Organization 1



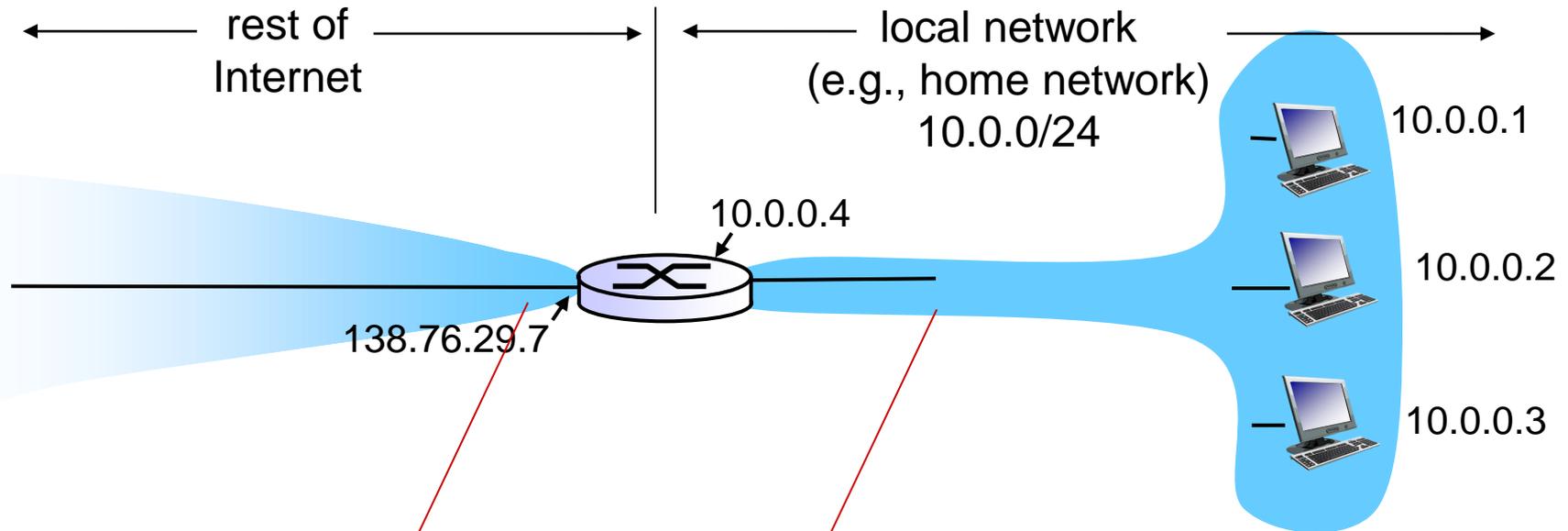
IP addressing: the last word...

Q: how does an ISP get block of addresses?

A: ICANN: Internet Corporation for Assigned Names and Numbers <http://www.icann.org/>

- allocates addresses
- manages DNS
- assigns domain names, resolves disputes

NAT: network address translation



all datagrams *leaving* local network have *same* single source NAT IP address: 138.76.29.7, different source port numbers

datagrams with source or destination in this network have 10.0.0/24 address for source, destination (as usual)

NAT: network address translation

motivation: local network uses just one IP address as far as outside world is concerned:

- range of addresses not needed from ISP: just one IP address for all devices
- can change addresses of devices in local network without notifying outside world
- can change ISP without changing addresses of devices in local network
- devices inside local net not explicitly addressable, visible by outside world (a security plus)

NAT: network address translation

implementation: NAT router must:

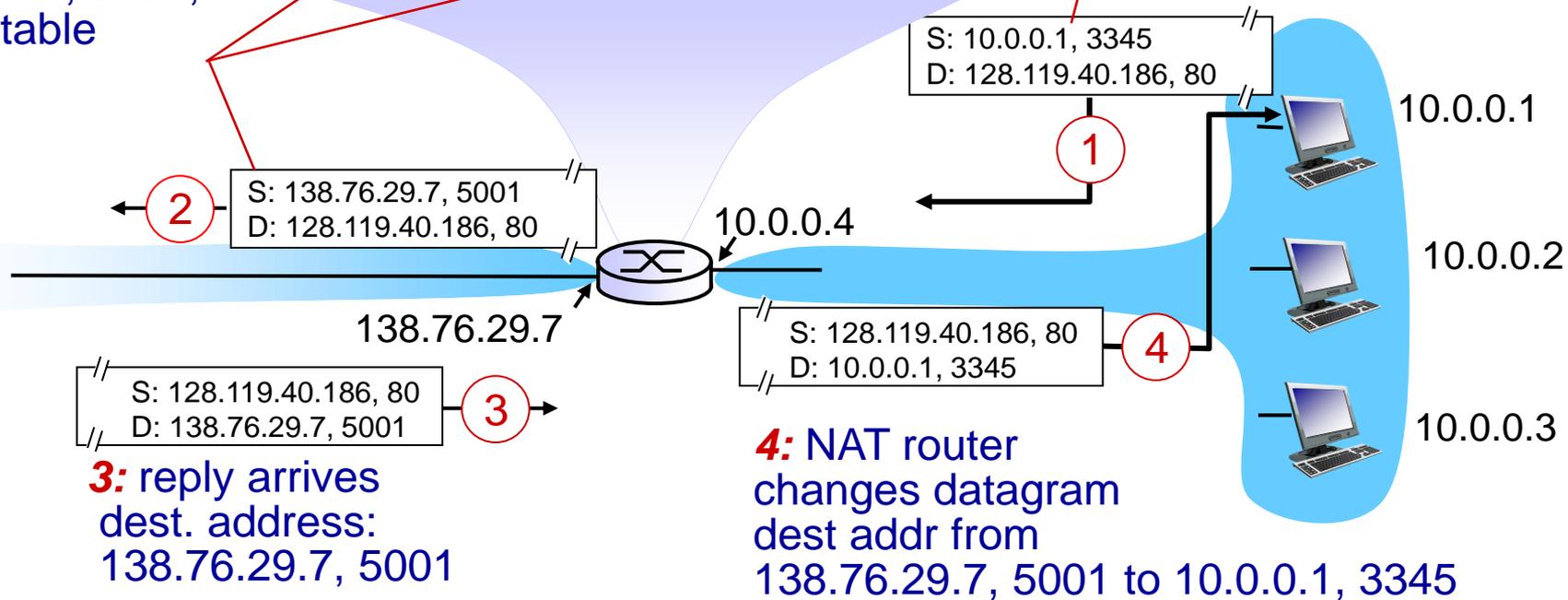
- *outgoing datagrams: replace* (source IP address, port #) of every outgoing datagram to (NAT IP address, new port #)
 - . . . remote clients/servers will respond using (NAT IP address, new port #) as destination addr
- *remember (in NAT translation table)* every (source IP address, port #) to (NAT IP address, new port #) translation pair
- *incoming datagrams: replace* (NAT IP address, new port #) in dest fields of every incoming datagram with corresponding (source IP address, port #) stored in NAT table

NAT: network address translation

NAT translation table	
WAN side addr	LAN side addr
138.76.29.7, 5001	10.0.0.1, 3345
.....

2: NAT router changes datagram source addr from 10.0.0.1, 3345 to 138.76.29.7, 5001, updates table

1: host 10.0.0.1 sends datagram to 128.119.40.186, 80

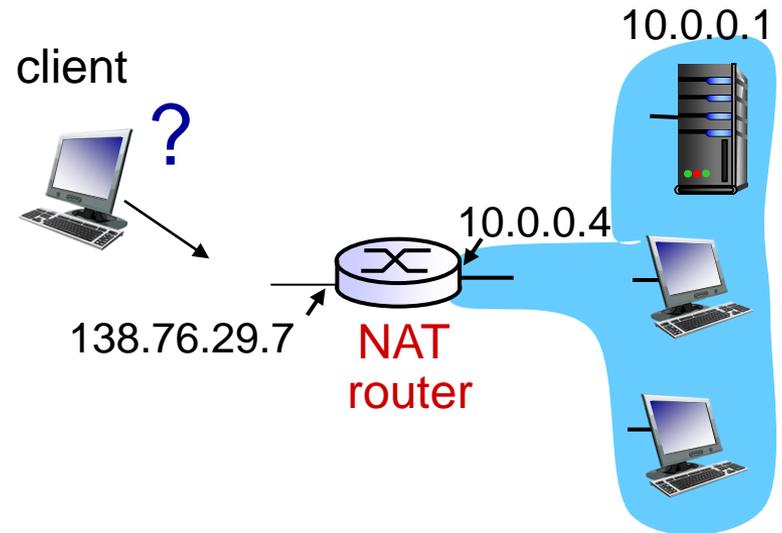


NAT: network address translation

- ❖ 16-bit port-number field:
 - 60,000 simultaneous connections with a single LAN-side address!
- ❖ NAT is controversial:
 - routers should only process up to layer 3
 - violates end-to-end argument
 - NAT possibility must be taken into account by app designers, e.g., P2P applications
 - address shortage should instead be solved by IPv6

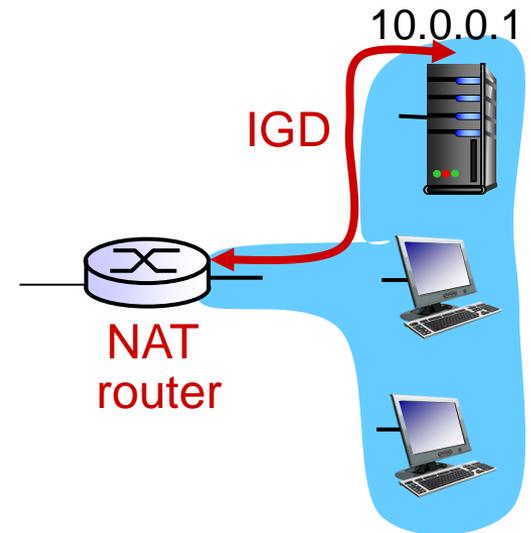
NAT traversal problem

- ❖ client wants to connect to server with address 10.0.0.1
 - server address 10.0.0.1 local to LAN (client can't use it as destination addr)
 - only one externally visible NATed address: 138.76.29.7
- ❖ **solution1:** statically configure NAT to forward incoming connection requests at given port to server
 - e.g., (123.76.29.7, port 2500) always forwarded to 10.0.0.1 port 25000



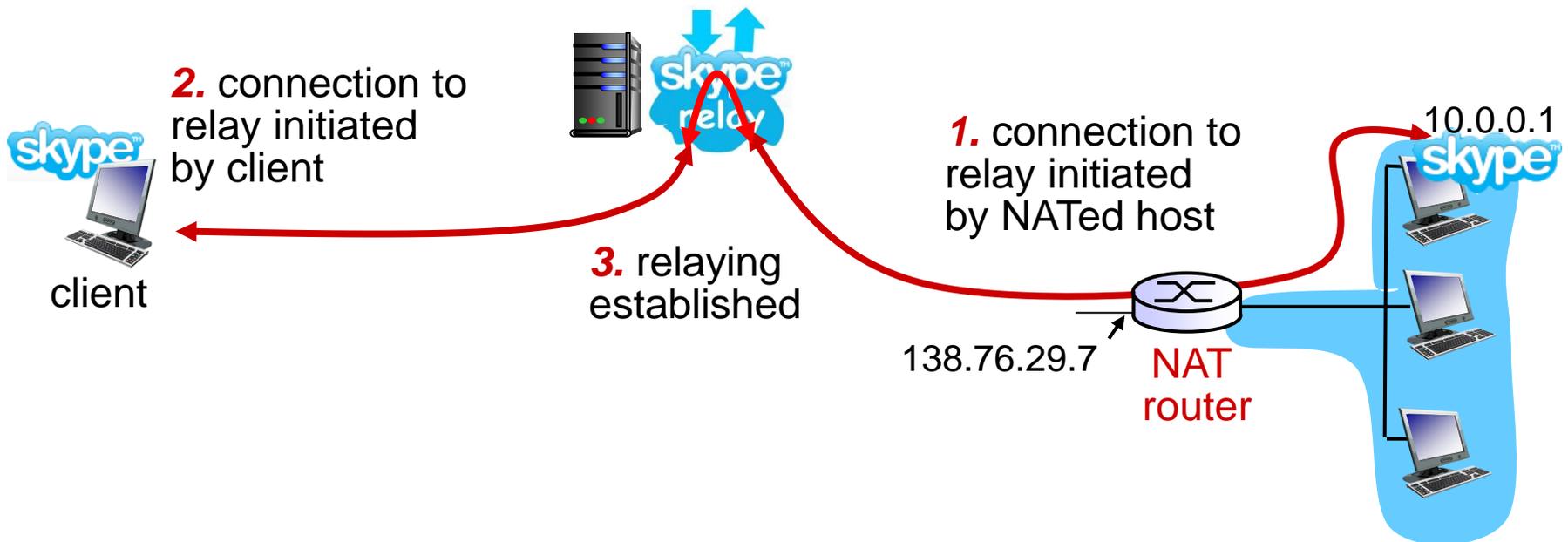
NAT traversal problem

- ❖ *solution 2*: Universal Plug and Play (UPnP) Internet Gateway Device (IGD) Protocol. Allows NATed host to:
 - ❖ learn public IP address (138.76.29.7)
 - ❖ add/remove port mappings (with lease times)
- i.e., automate static NAT port map configuration



NAT traversal problem

- ❖ **solution 3:** relaying (used in Skype)
 - NATed client establishes connection to relay
 - external client connects to relay
 - relay bridges packets between to connections



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4.3 what's inside a router

4.4 IP: Internet Protocol

- datagram format
- IPv4 addressing
- ICMP
- IPv6

4.5 routing algorithms

- link state
- distance vector
- hierarchical routing

4.6 routing in the Internet

- RIP
- OSPF
- BGP

4.7 broadcast and multicast routing

ICMP: internet control message protocol

- ❖ used by hosts & routers to communicate network-level information
 - error reporting: unreachable host, network, port, protocol
 - echo request/reply (used by ping)
- ❖ network-layer “above” IP:
 - ICMP msgs carried in IP datagrams
- ❖ **ICMP message:** type, code plus first 8 bytes of IP datagram causing error

<u>Type</u>	<u>Code</u>	<u>description</u>
0	0	echo reply (ping)
3	0	dest. network unreachable
3	1	dest host unreachable
3	2	dest protocol unreachable
3	3	dest port unreachable
3	6	dest network unknown
3	7	dest host unknown
4	0	source quench (congestion control - not used)
8	0	echo request (ping)
9	0	route advertisement
10	0	router discovery
11	0	TTL expired
12	0	bad IP header

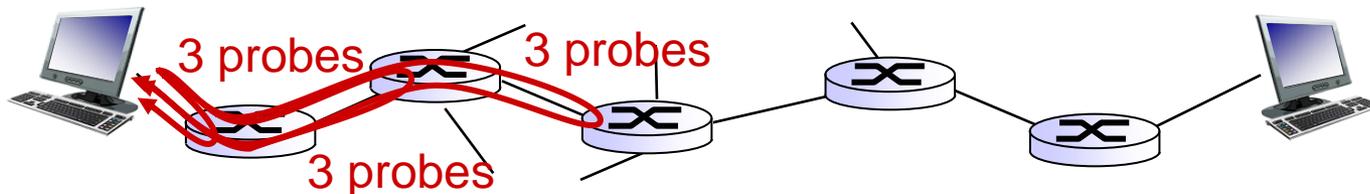
Traceroute and ICMP

- ❖ source sends series of UDP segments to dest
 - first set has TTL = 1
 - second set has TTL=2, etc.
 - unlikely port number
- ❖ when n th set of datagrams arrives to n th router:
 - router discards datagrams
 - and sends source ICMP messages (type 11, code 0)
 - ICMP messages includes name of router & IP address

- ❖ when ICMP messages arrives, source records RTTs

stopping criteria:

- ❖ UDP segment eventually arrives at destination host
- ❖ destination returns ICMP “port unreachable” message (type 3, code 3)
- ❖ source stops



IPv6: motivation

- ❖ *initial motivation*: 32-bit address space soon to be completely allocated.
- ❖ additional motivation:
 - header format helps speed processing/forwarding
 - header changes to facilitate QoS

IPv6 datagram format:

- fixed-length 40 byte header
- no fragmentation allowed

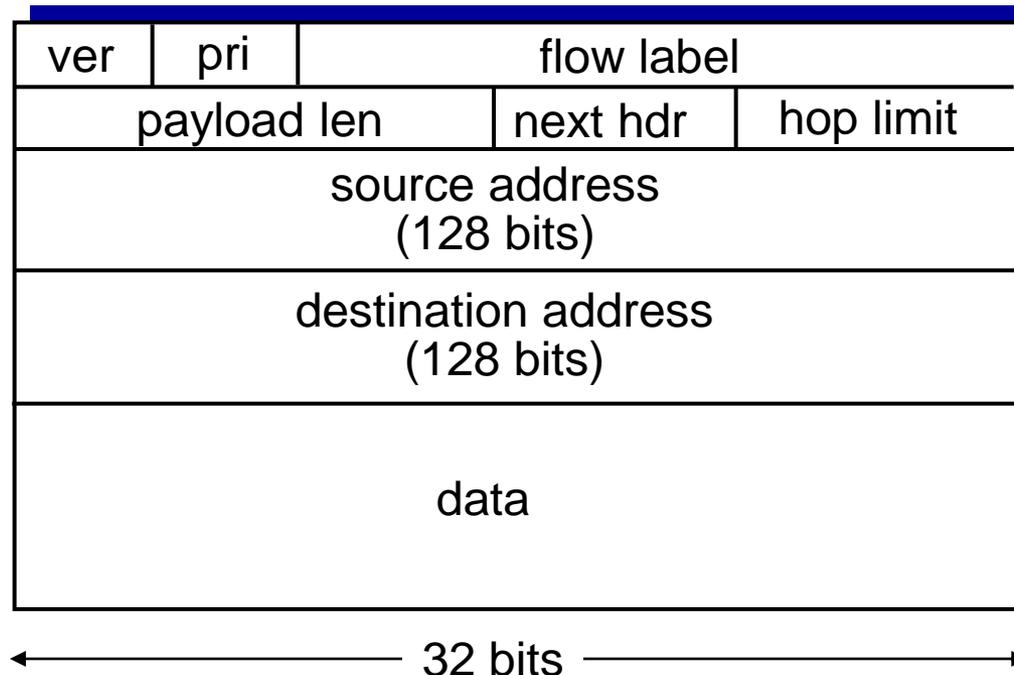
IPv6 datagram format

priority: identify priority among datagrams in flow

flow Label: identify datagrams in same “flow.”

(concept of “flow” not well defined).

next header: identify upper layer protocol for data

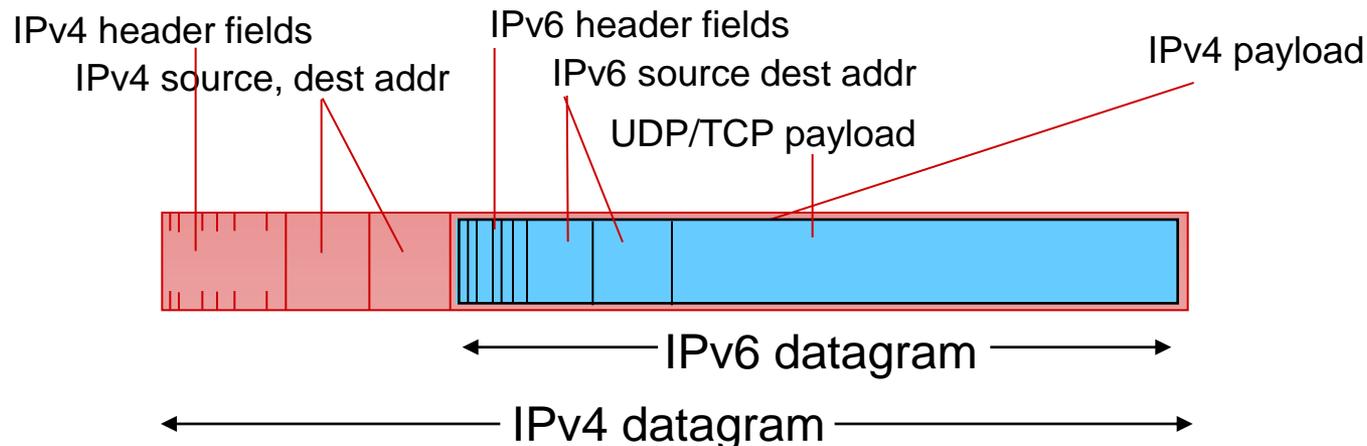


Other changes from IPv4

- ❖ *checksum*: removed entirely to reduce processing time at each hop
- ❖ *options*: allowed, but outside of header, indicated by “Next Header” field
- ❖ *ICMPv6*: new version of ICMP
 - additional message types, e.g. “Packet Too Big”
 - multicast group management functions

Transition from IPv4 to IPv6

- ❖ not all routers can be upgraded simultaneously
 - no “flag days”
 - how will network operate with mixed IPv4 and IPv6 routers?
- ❖ *tunneling*: IPv6 datagram carried as *payload* in IPv4 datagram among IPv4 routers

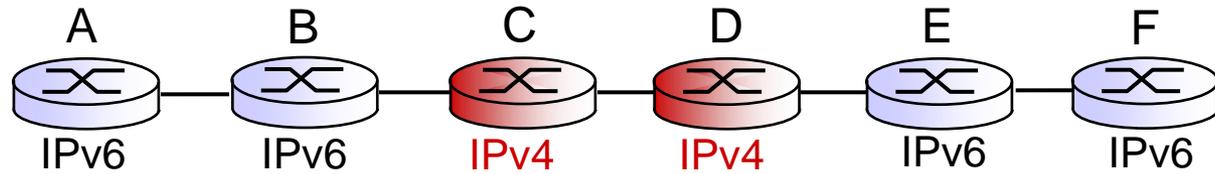


Tunneling

logical view:

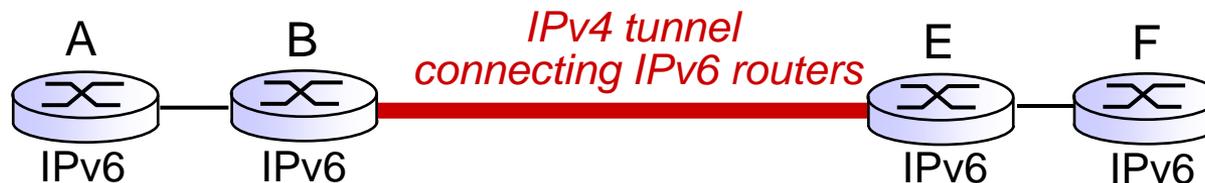


physical view:



Tunneling

logical view:



physical view:

