

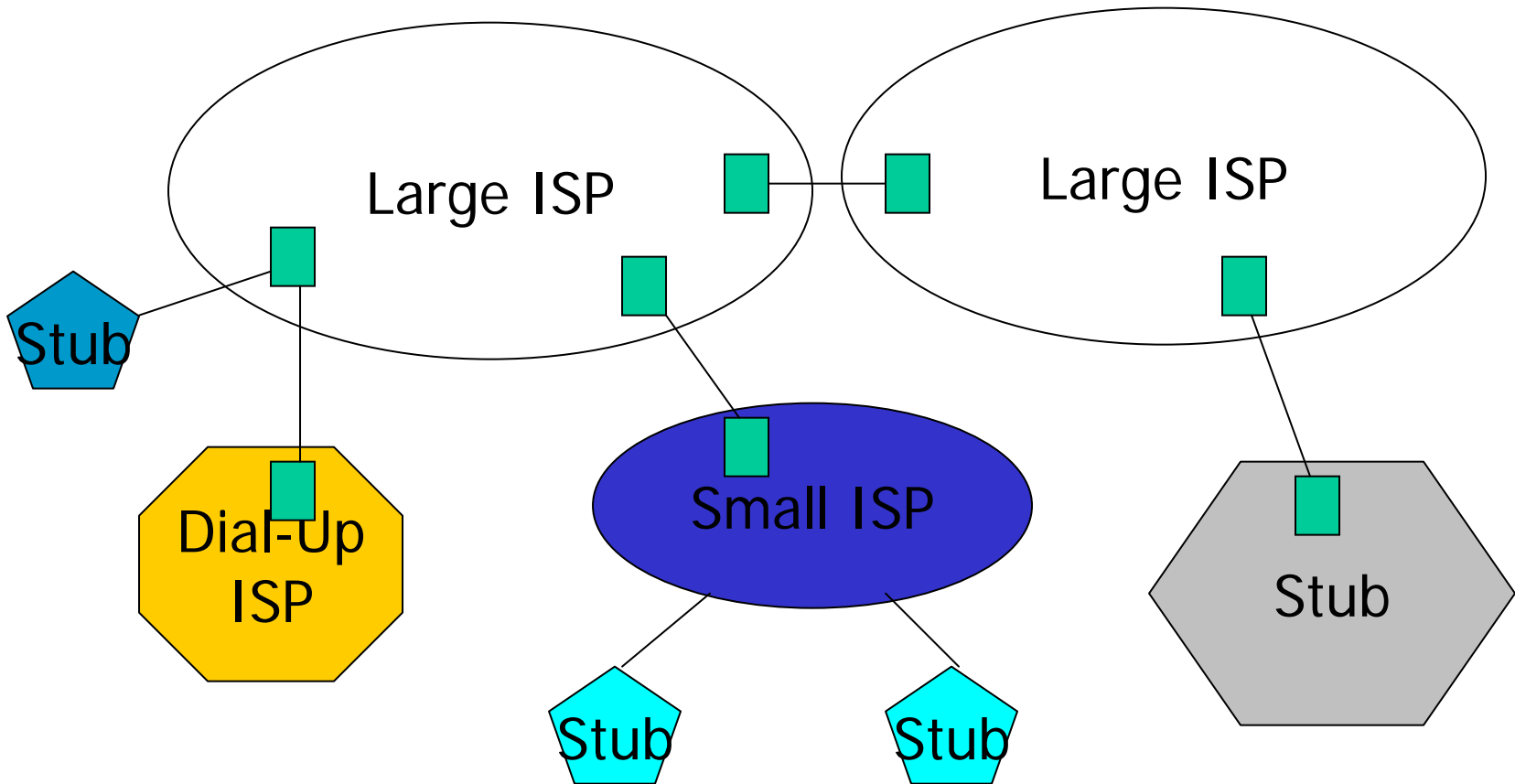
# COMP/ELEC 429

## Introduction to Computer Networks

Lecture 11: Inter-domain routing

Slides used with permissions from Edward W. Knightly,  
T. S. Eugene Ng, Ion Stoica, Hui Zhang

# Internet Structure



# Autonomous Systems (AS)

- Internet is not a single network!
- The Internet is a collection of networks, each controlled by different administrations
- An autonomous system (AS) is a network under a single administrative control

# AS Numbers (ASNs)

ASNs are 16 bit values.

64512 through 65535 are “private”

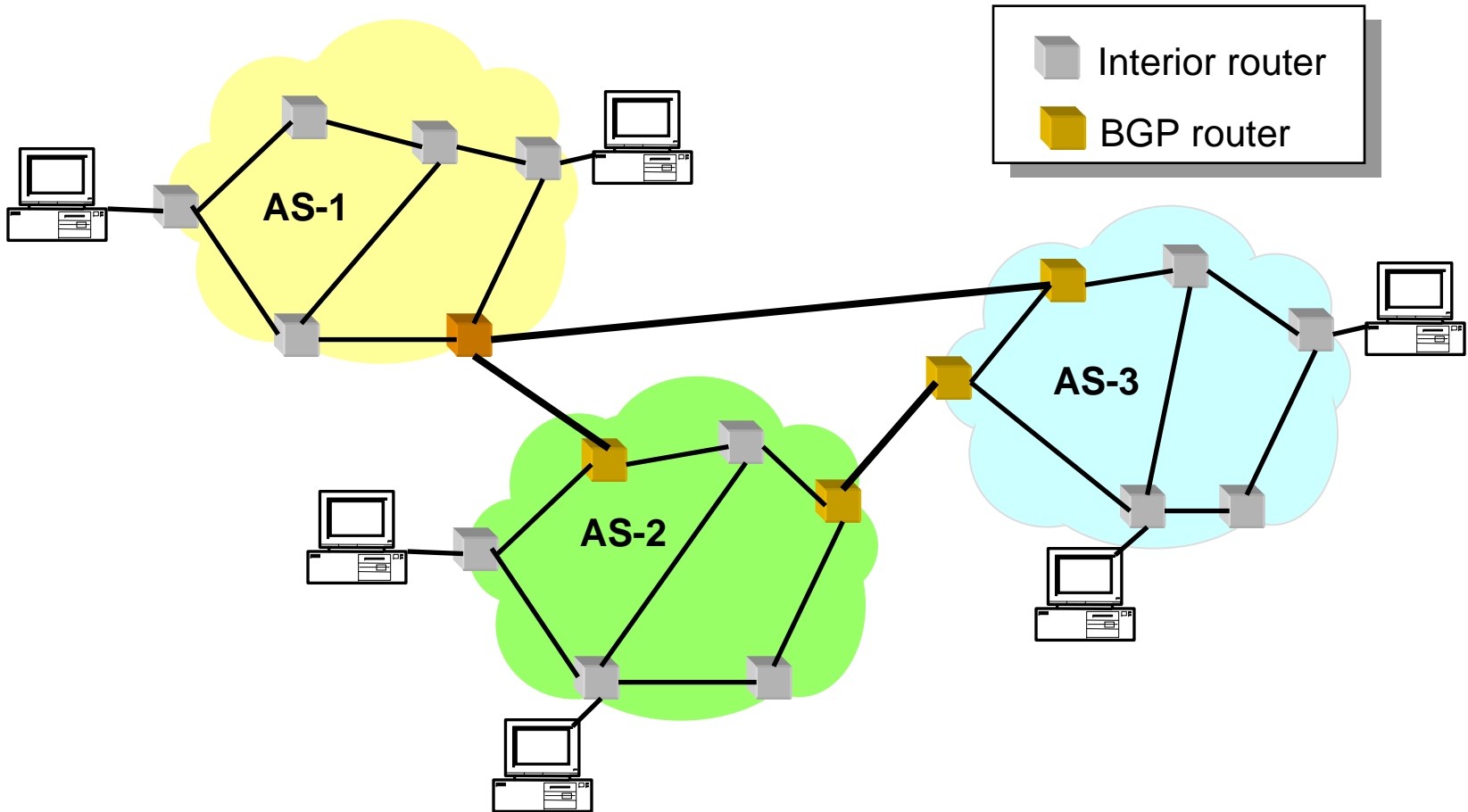
Currently over 11,000 in use.

- Genuity: 1
- AT&T: 7018, 6341, 5074, ...
- UUNET: 701, 702, 284, 12199, ...
- Sprint: 1239, 1240, 6211, 6242, ...
- ...

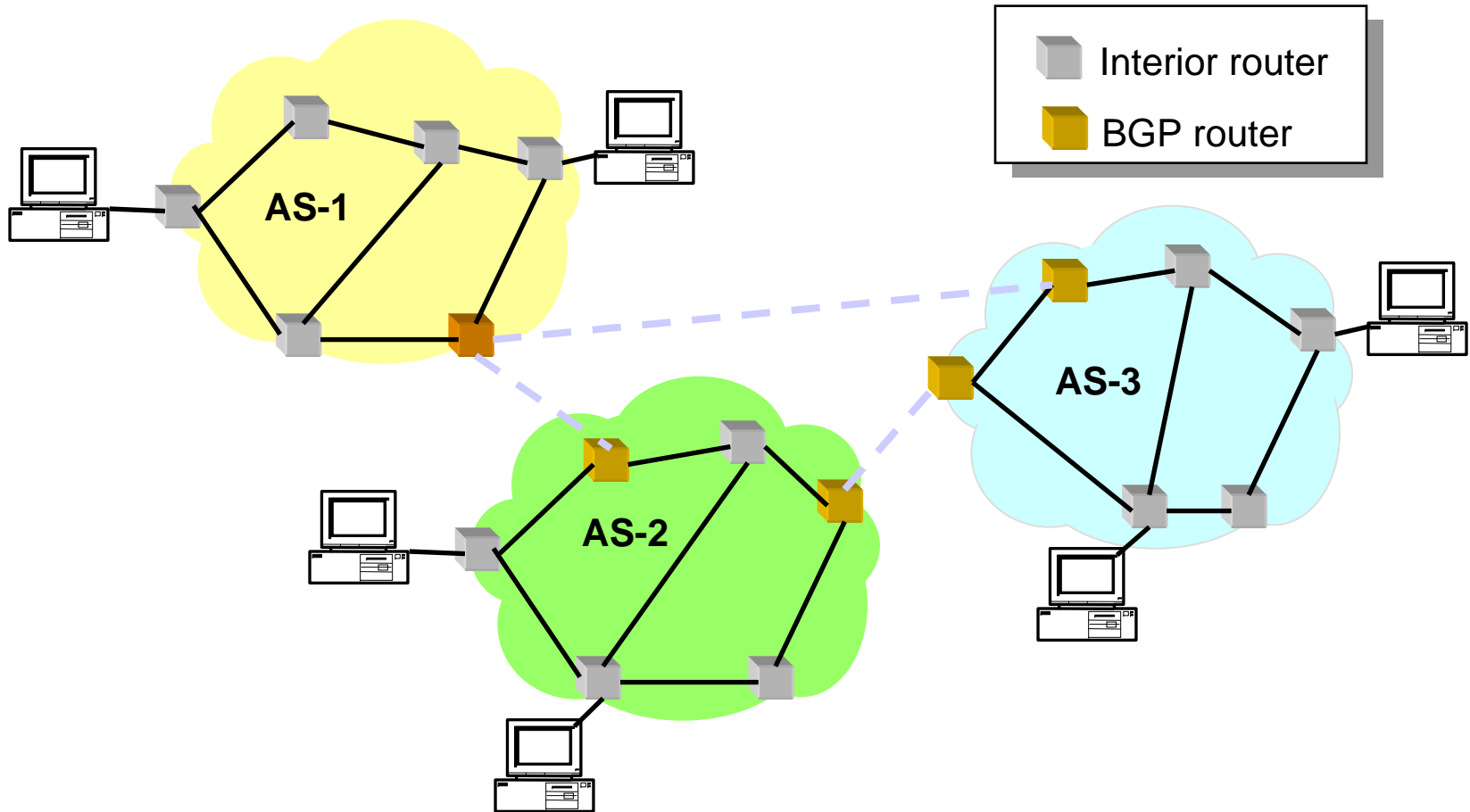
# Implications

- ASs want to choose own local routing algorithm
  - AS takes care of getting packets to/from their own hosts
  - Intradomain routing: RIP, OSPF, etc
- ASs want to choose own non-local routing policy
  - Interdomain routing must accommodate this
  - BGP is the current interdomain routing protocol
  - BGP: Border Gateway Protocol

# Example

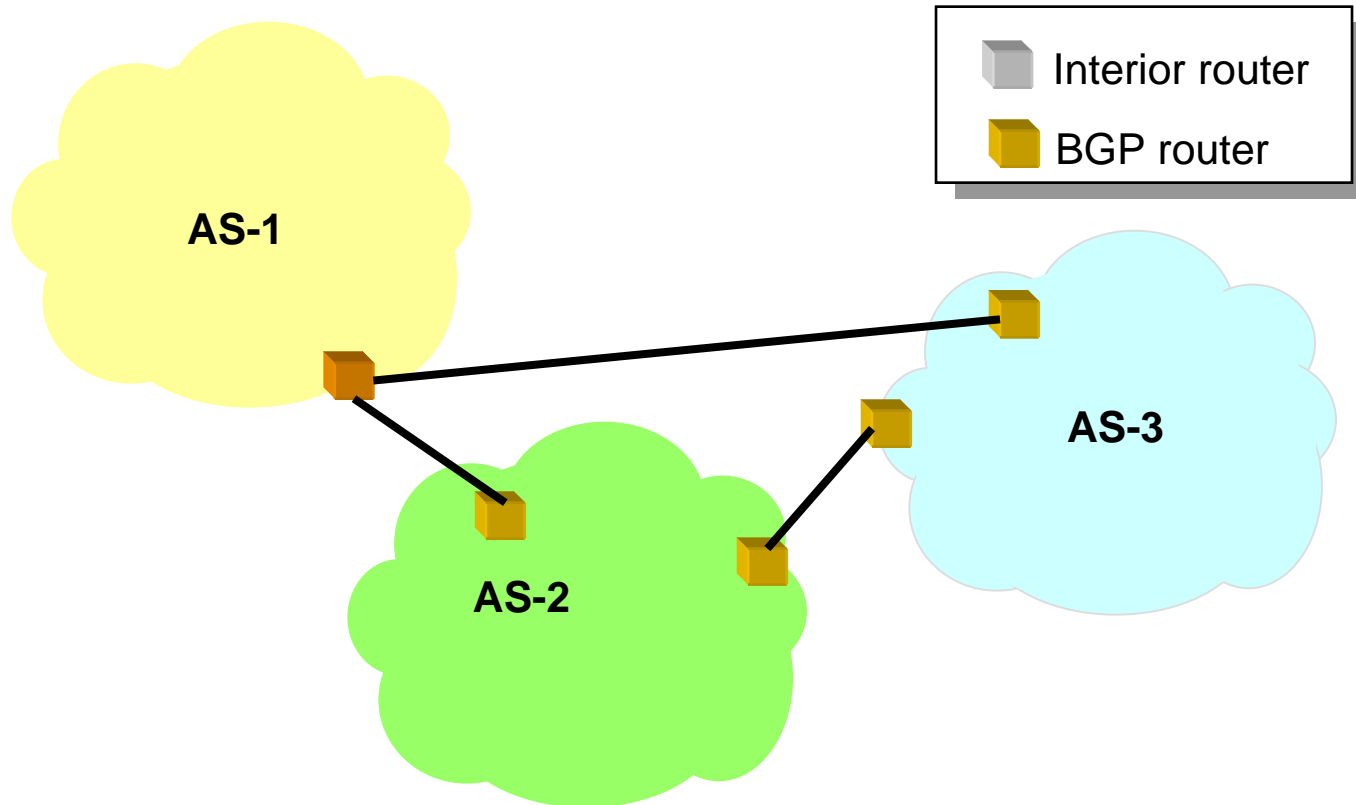


# Intra-Domain



Intra-domain routing protocol aka **Interior Gateway Protocol (IGP)**, e.g. OSPF, RIP

# Inter-Domain



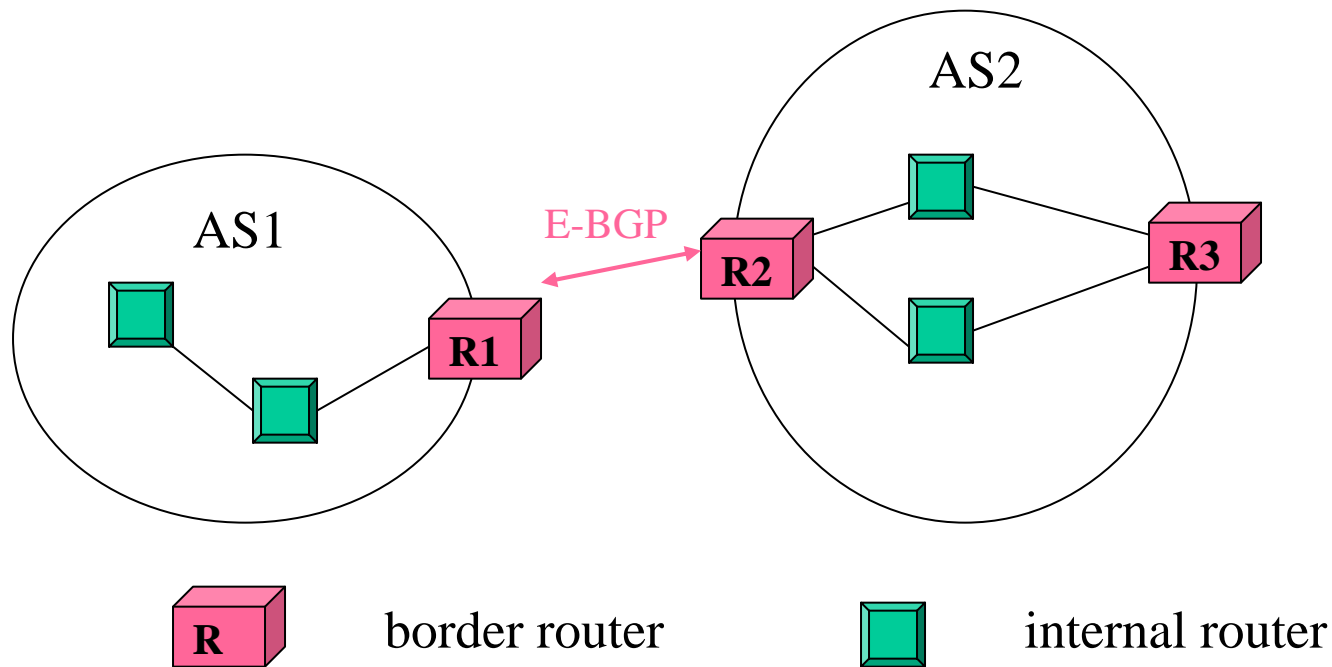
Inter-domain routing protocol aka **Exterior Gateway Protocol (EGP)**, e.g. BGP



# Inter-Domain Routing

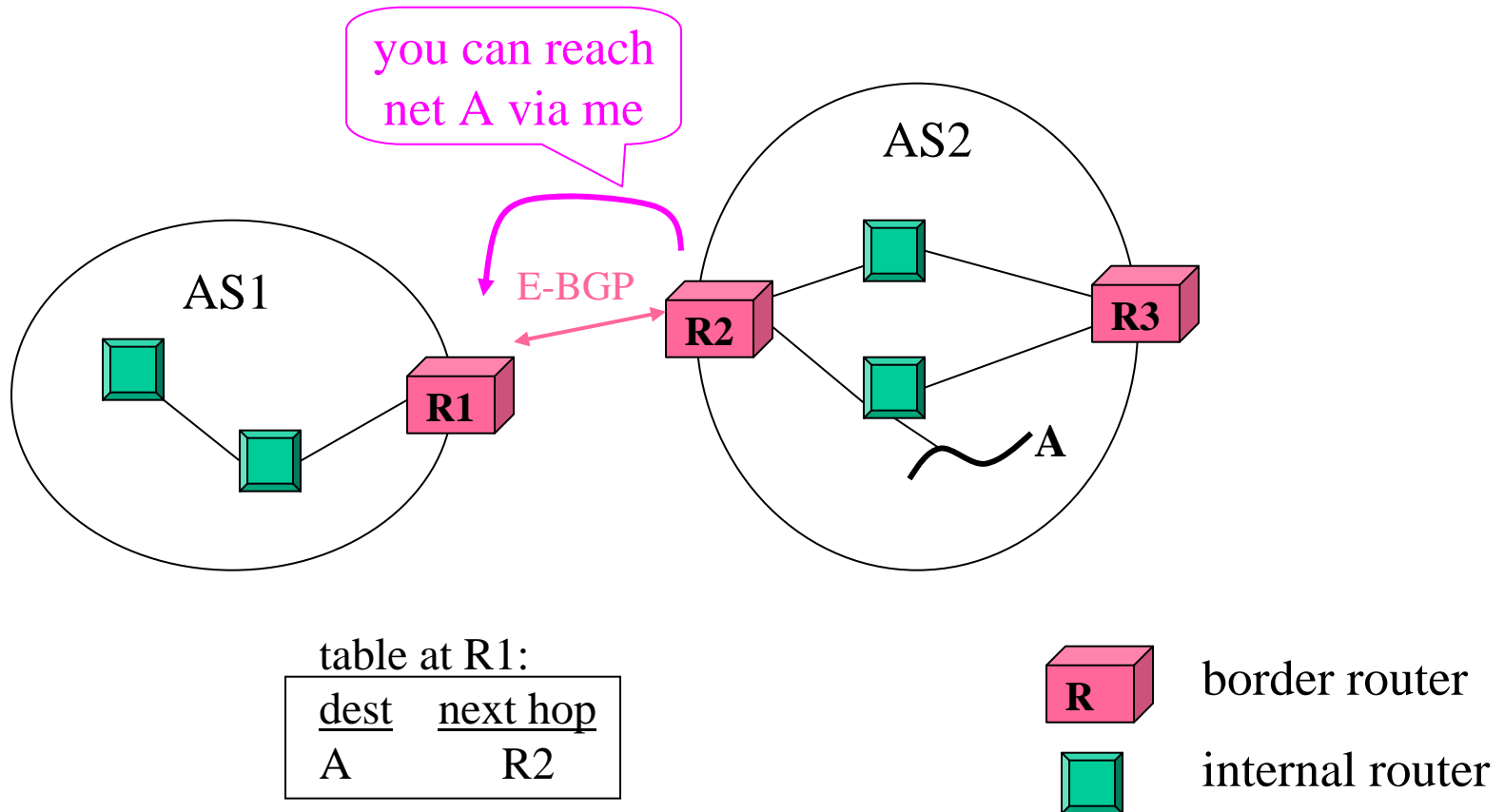
- Global connectivity is at stake
- Inevitably leads to one single protocol that everyone must speak
  - Unlike many choices in intra-domain routing
- What are the requirements?
- Scalability
- Flexibility in choosing routes
- If you were to choose, link state based or distance vector based?
  
- BGP is sort of a hybrid: Path vector protocol

# Border Gateway Protocol Part I: E-BGP



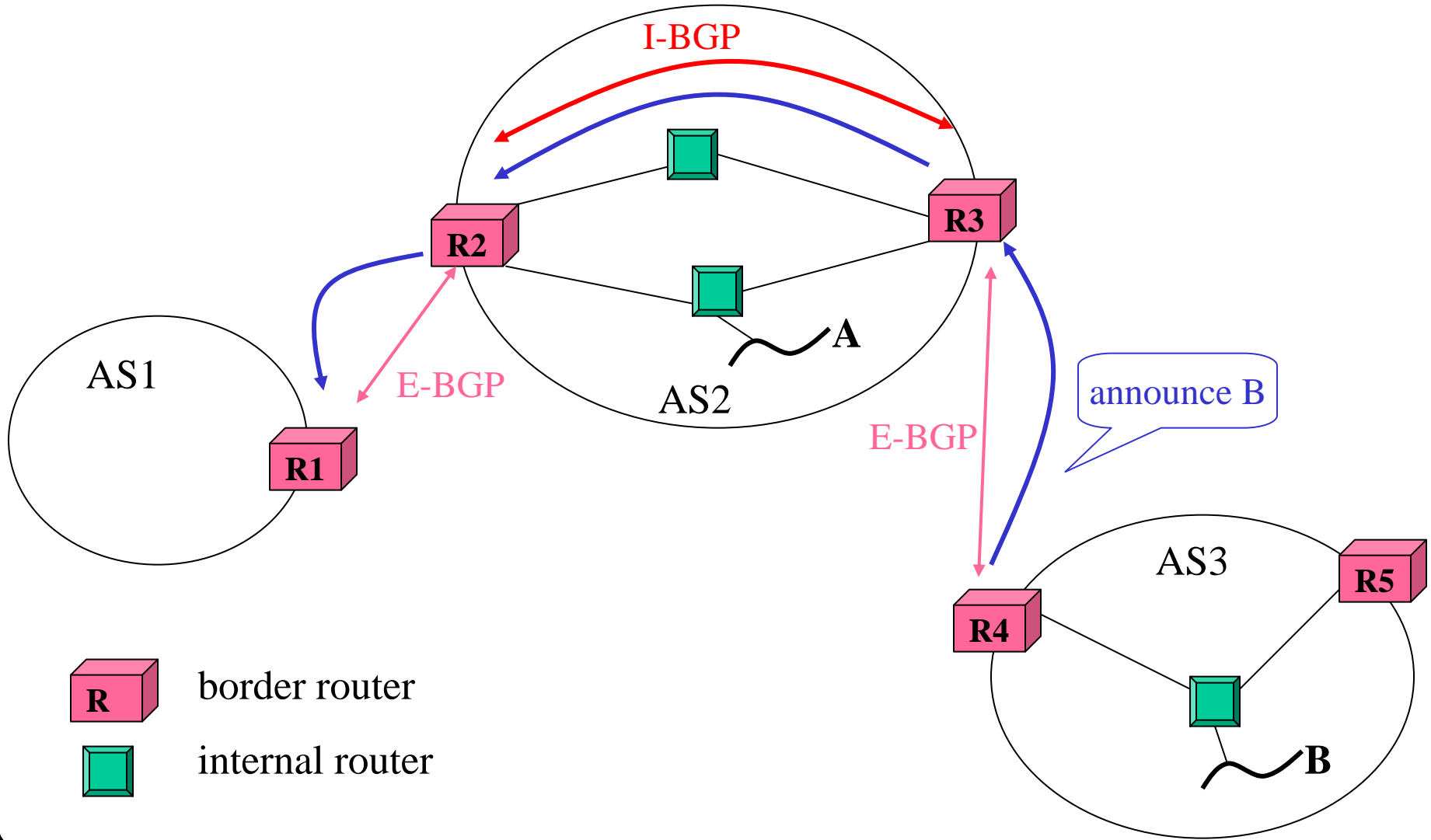
- Two types of routers
  - Border router (Edge), Internal router (Core)

# Purpose of E-BGP

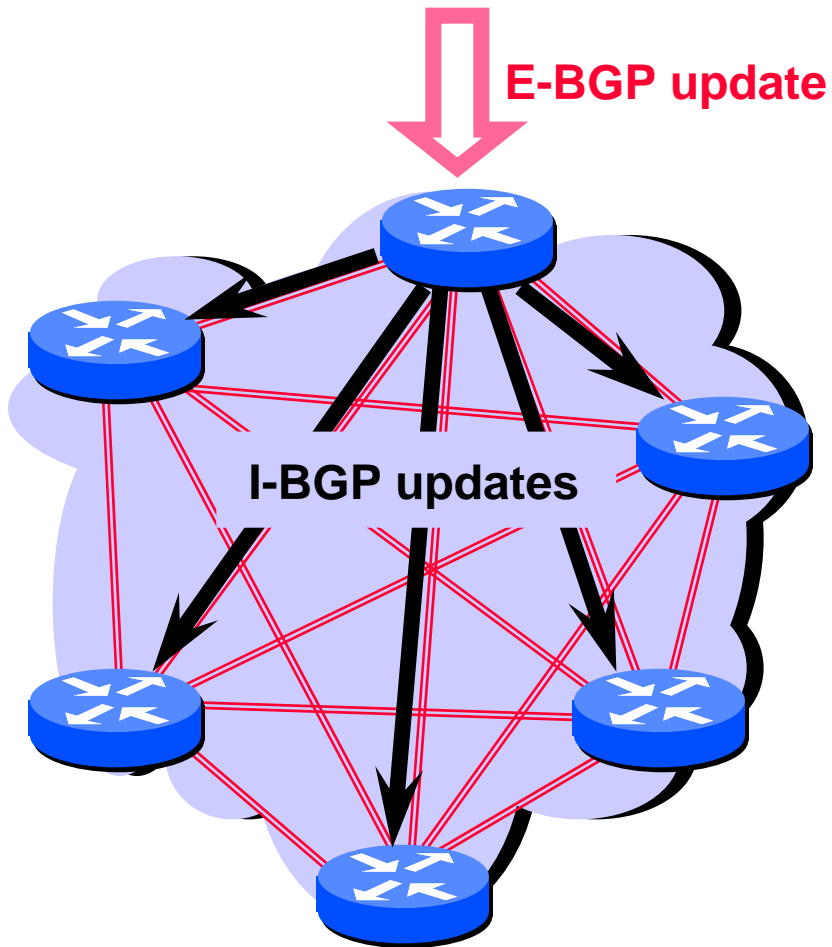


Share connectivity information across ASes

# Part II: I-BGP, Carrying Info within an AS



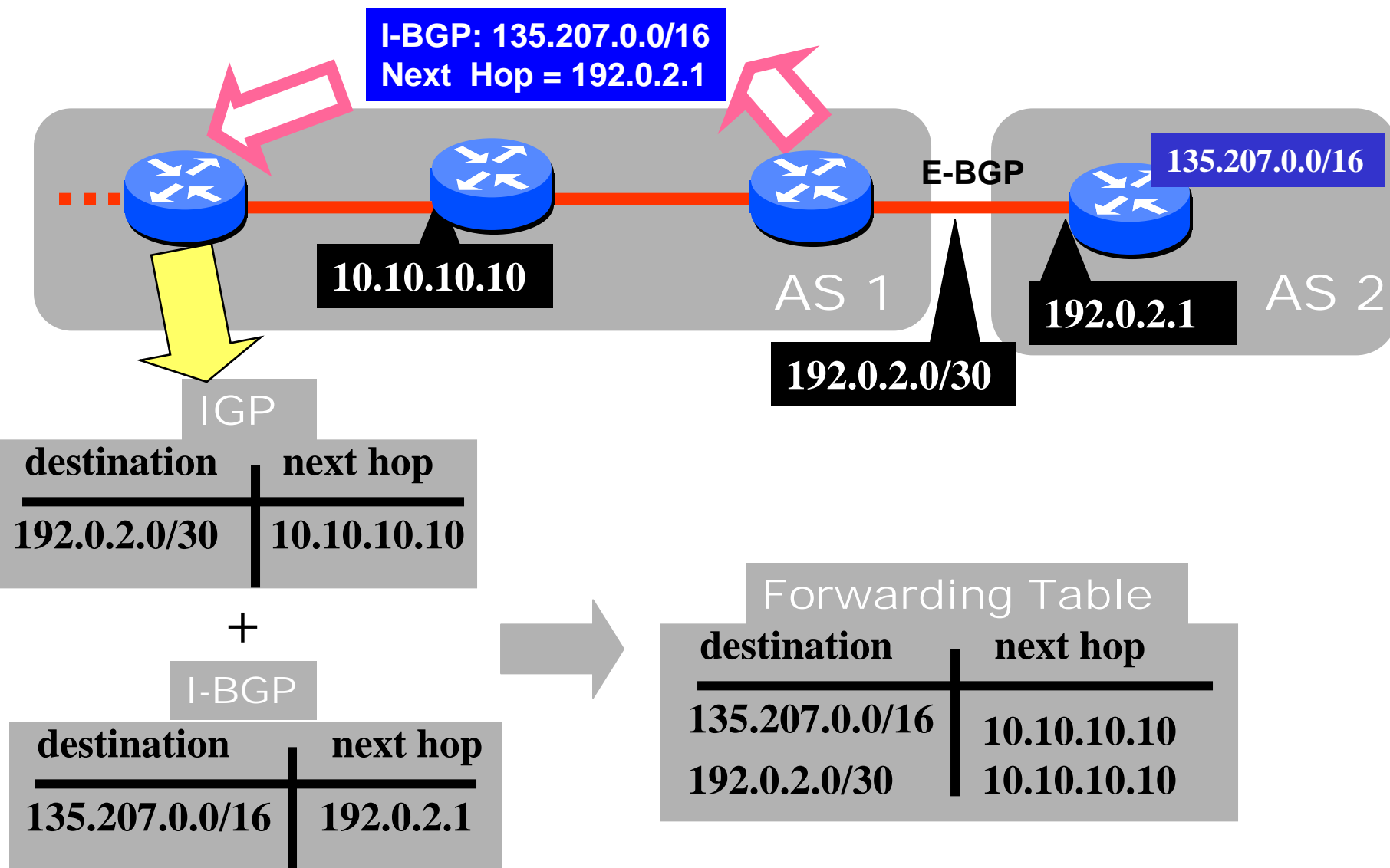
# I-BGP



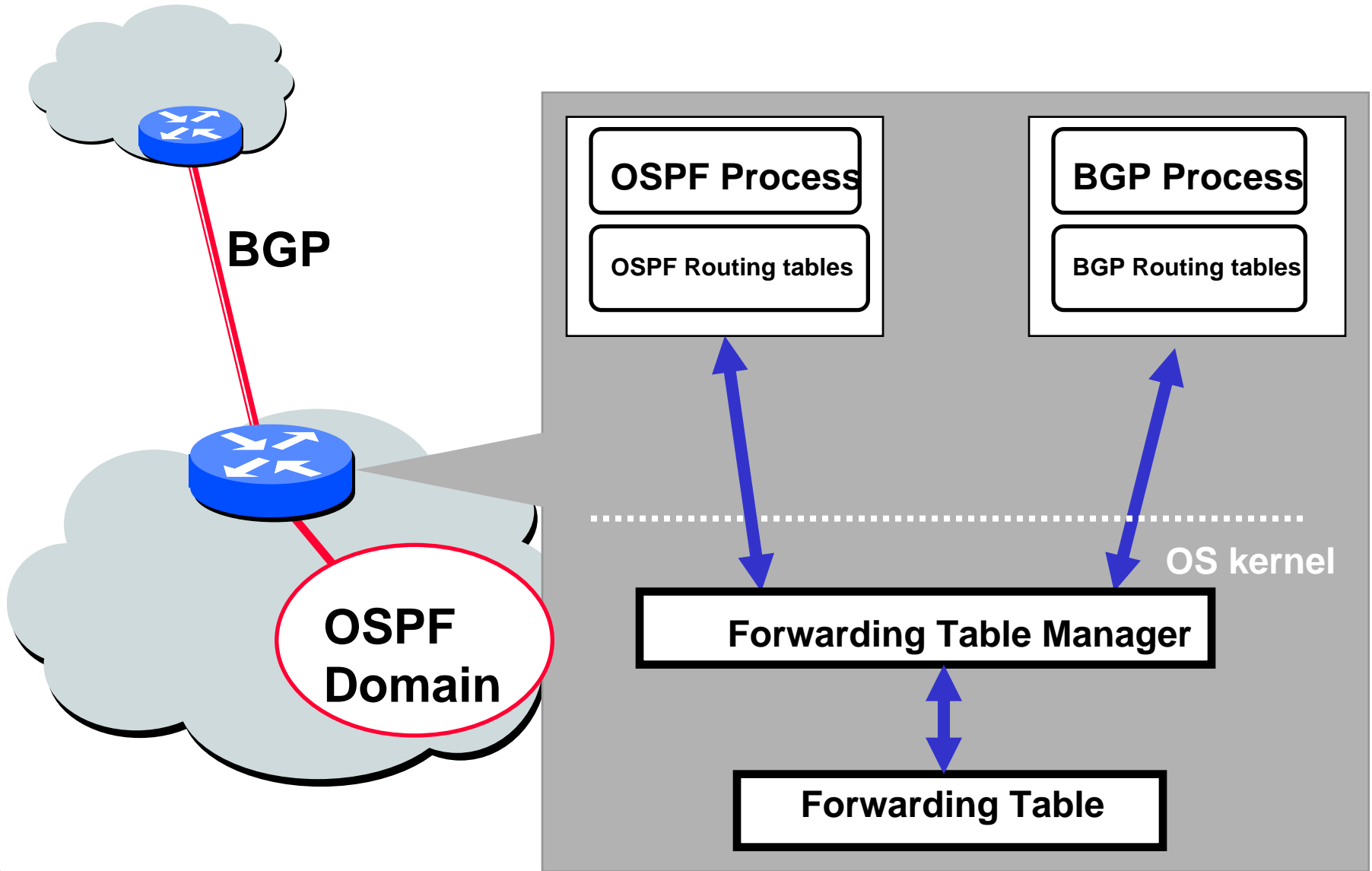
- Problem: Injecting external routes into IGP (e.g. OSPF) does not scale and causes BGP policy information to be lost
- I-BGP can be used to disseminate BGP routes to all routers in AS
- BGP route + IGP route suffice to create forwarding table

I-BGP neighbors do not announce routes received via I-BGP to other I-BGP neighbors.

# Join I-BGP with IGP to Create Forwarding Table



# Multiple Routing Processes on a Single Router

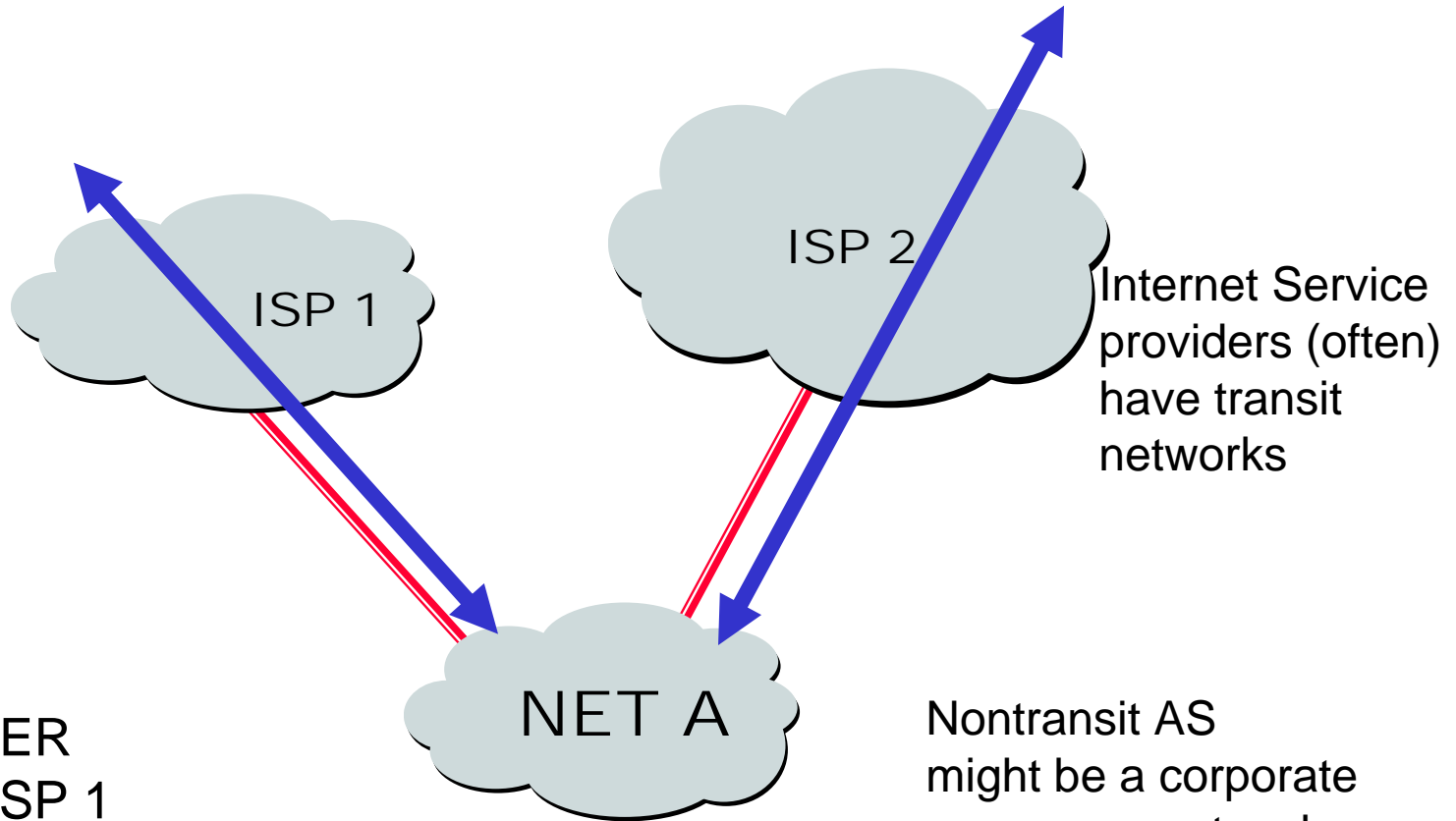


# Routing between ISPs

- Routing protocol (BGP) contains reachability information (no metrics)
  - Not about optimizing anything
  - All about policy (business and politics)
- Why?
  - Metrics optimize for a particular criteria
  - AT&T's idea of a good route is not the same as UUnet's
  - Scale
- What a BGP speaker announces or not announces to a peer determines what routes may get used by whom

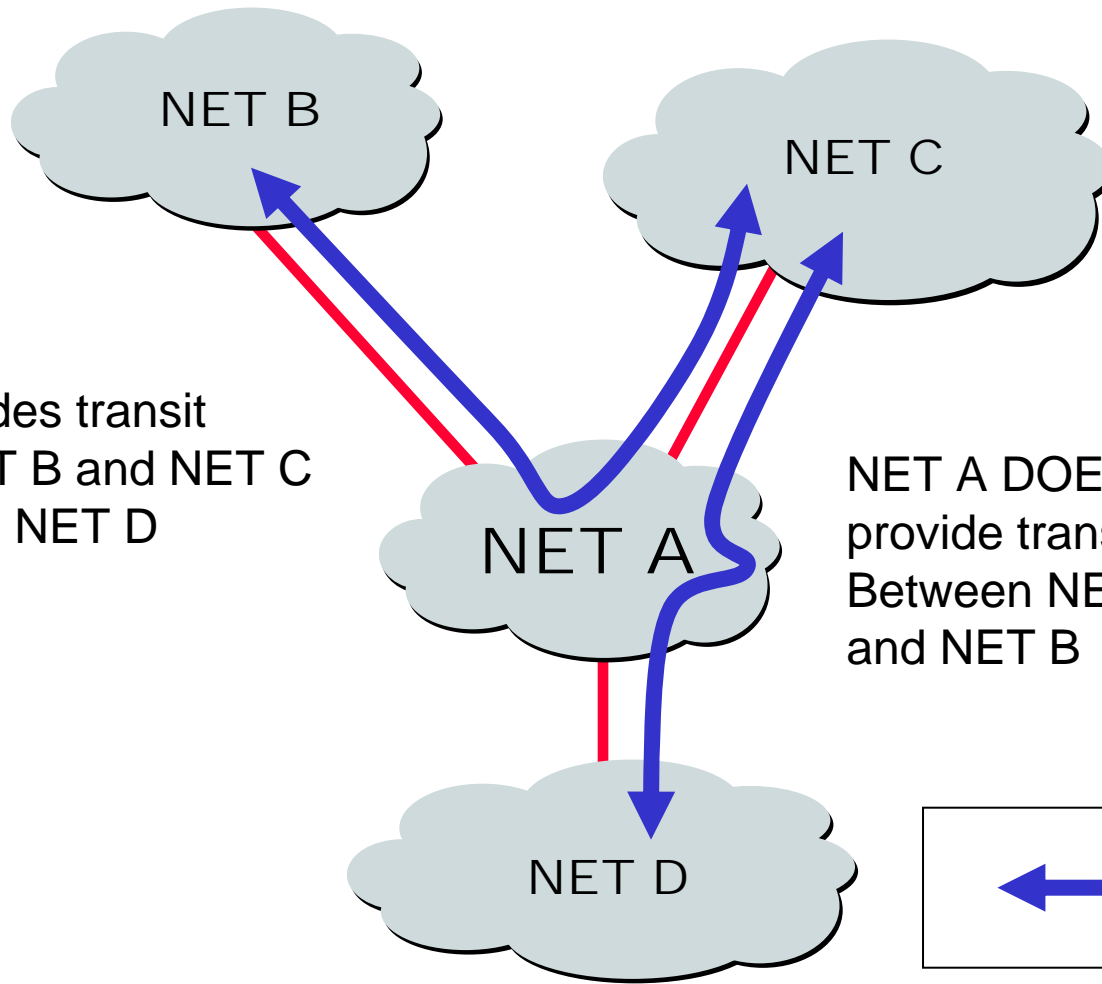


# Nontransit vs. Transit ASes



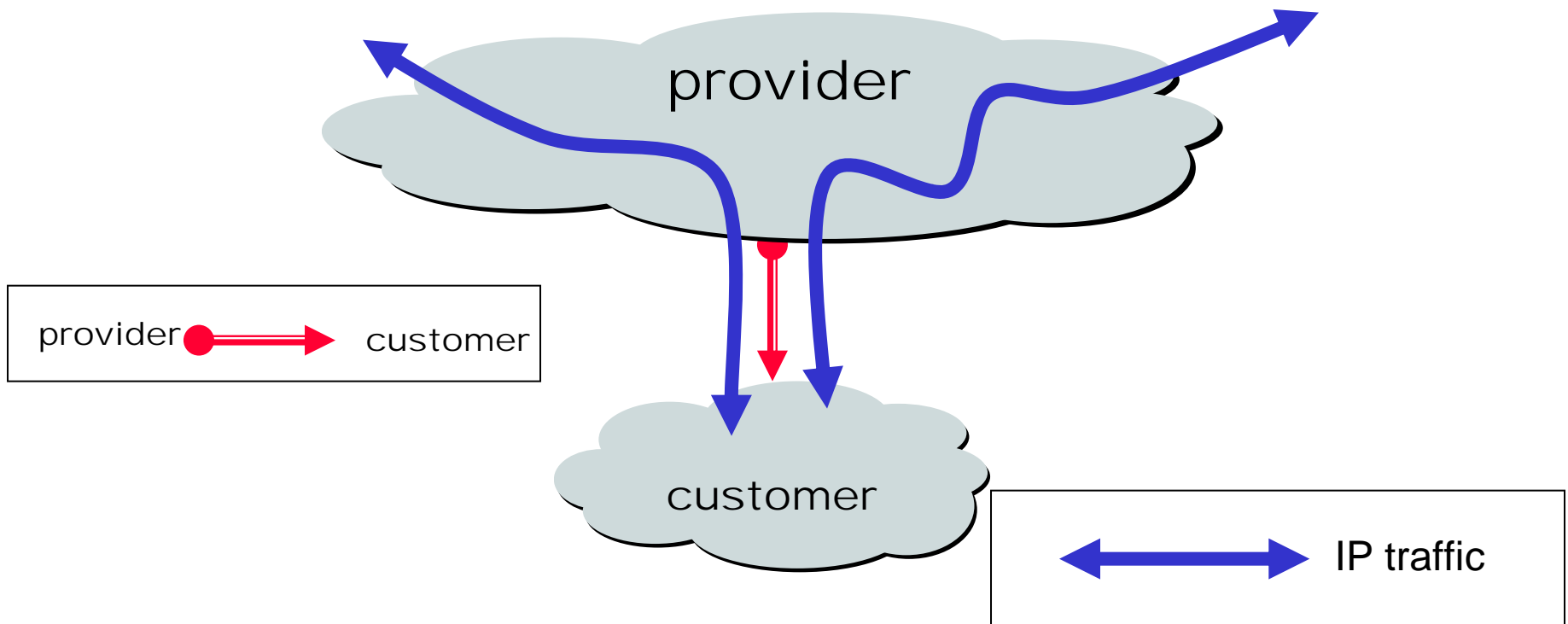
Traffic NEVER flows from ISP 1 through NET A to ISP 2 (At least not intentionally!)

# Selective Transit



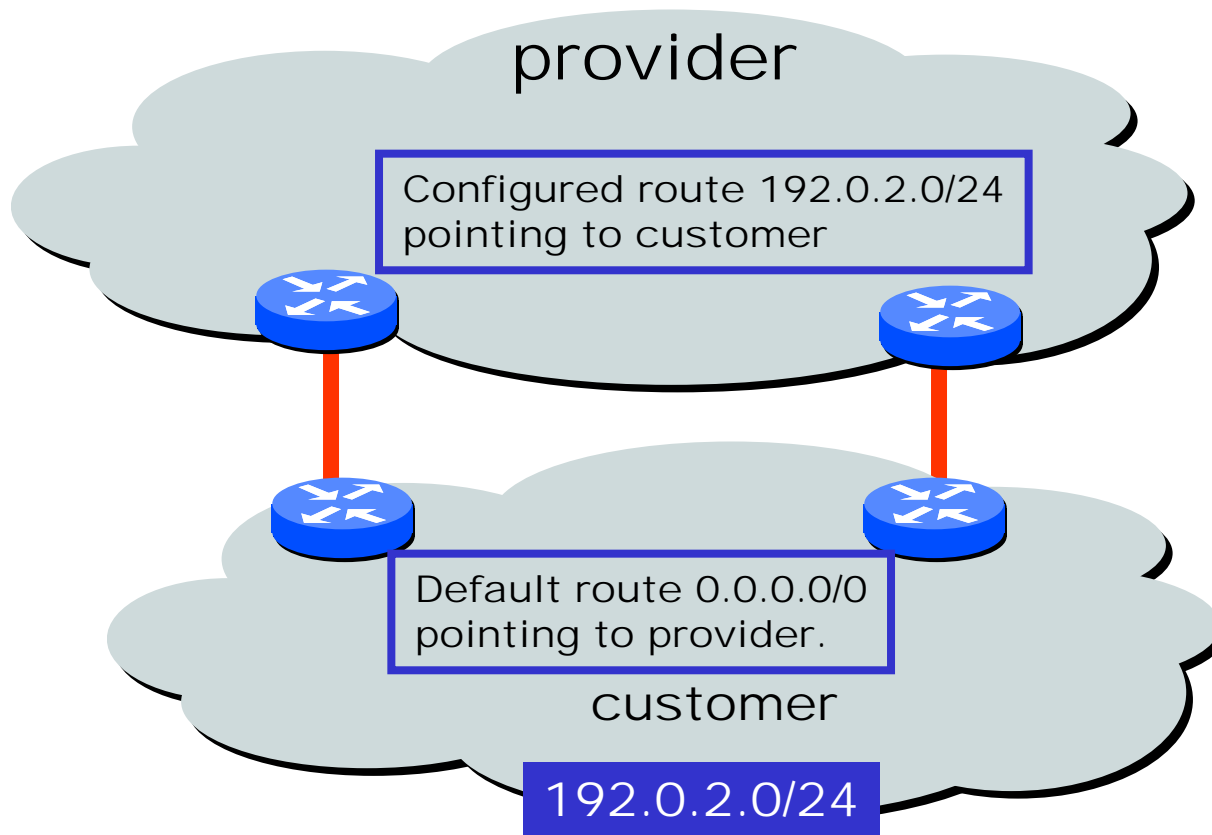
Most transit networks transit in a selective manner...

# Customers and Providers



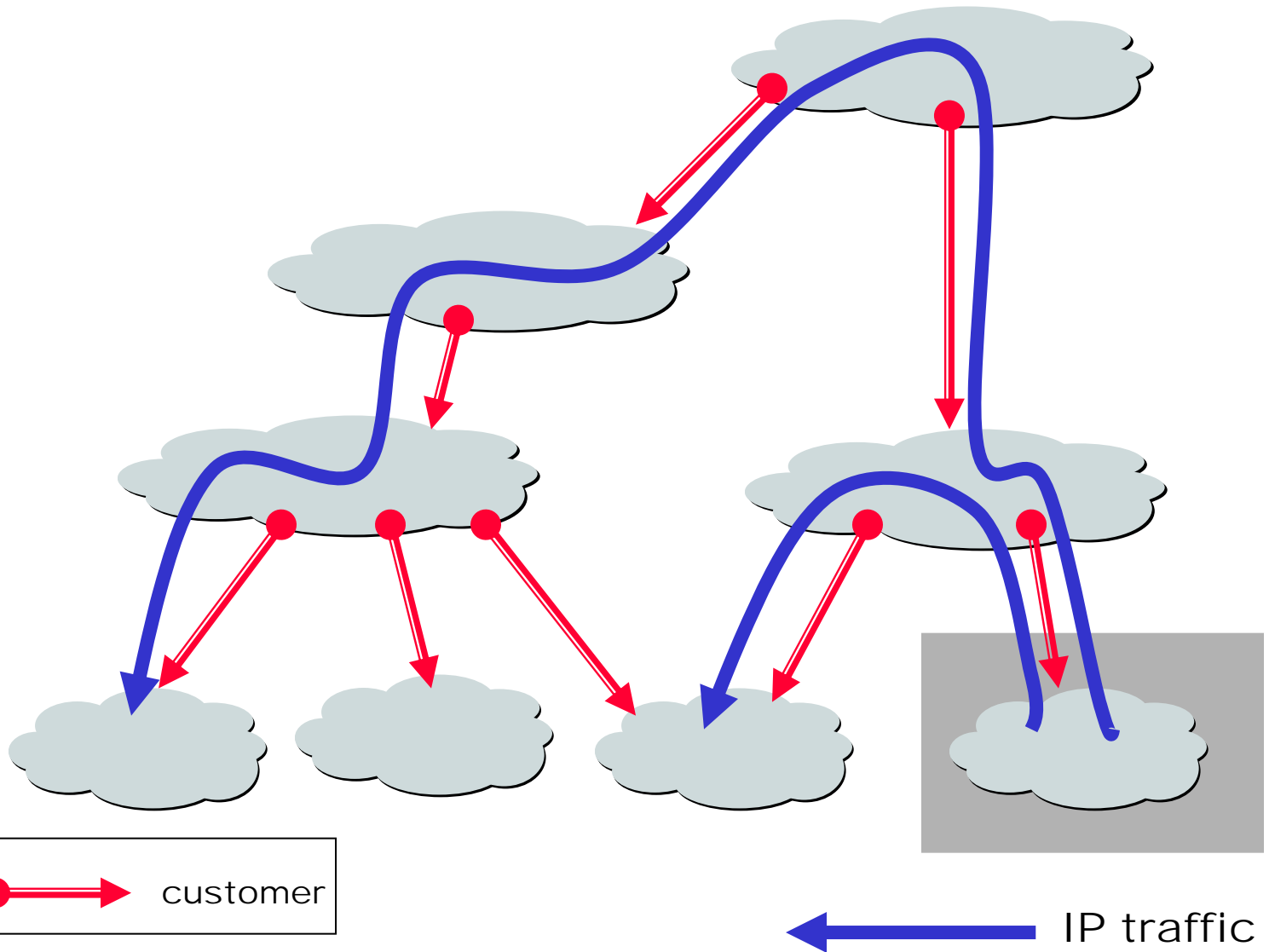
Customer pays provider for access to the Internet

# Customers Don't Always Need BGP

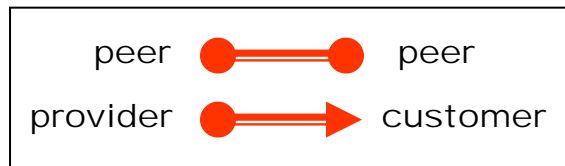
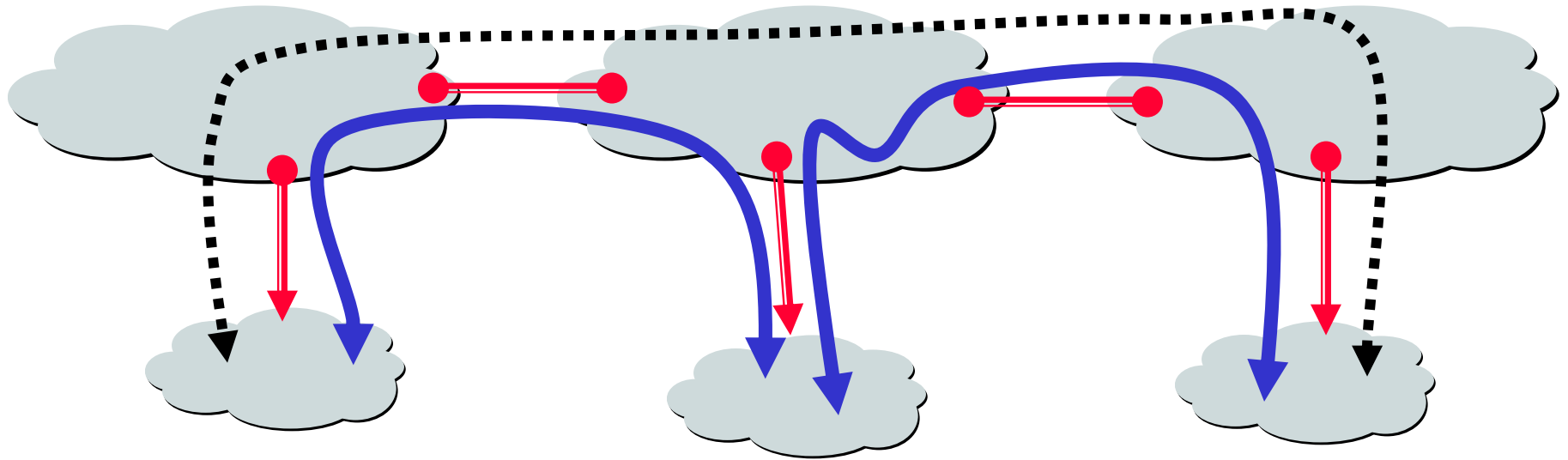


Static routing is the most common way of connecting an autonomous routing domain to the Internet. This helps explain why BGP is a mystery to many ...

# Customer-Provider Hierarchy



# The Peering Relationship



Peers provide transit between their respective customers

Peers do not provide transit between peers

Peers (often) do not exchange \$\$\$

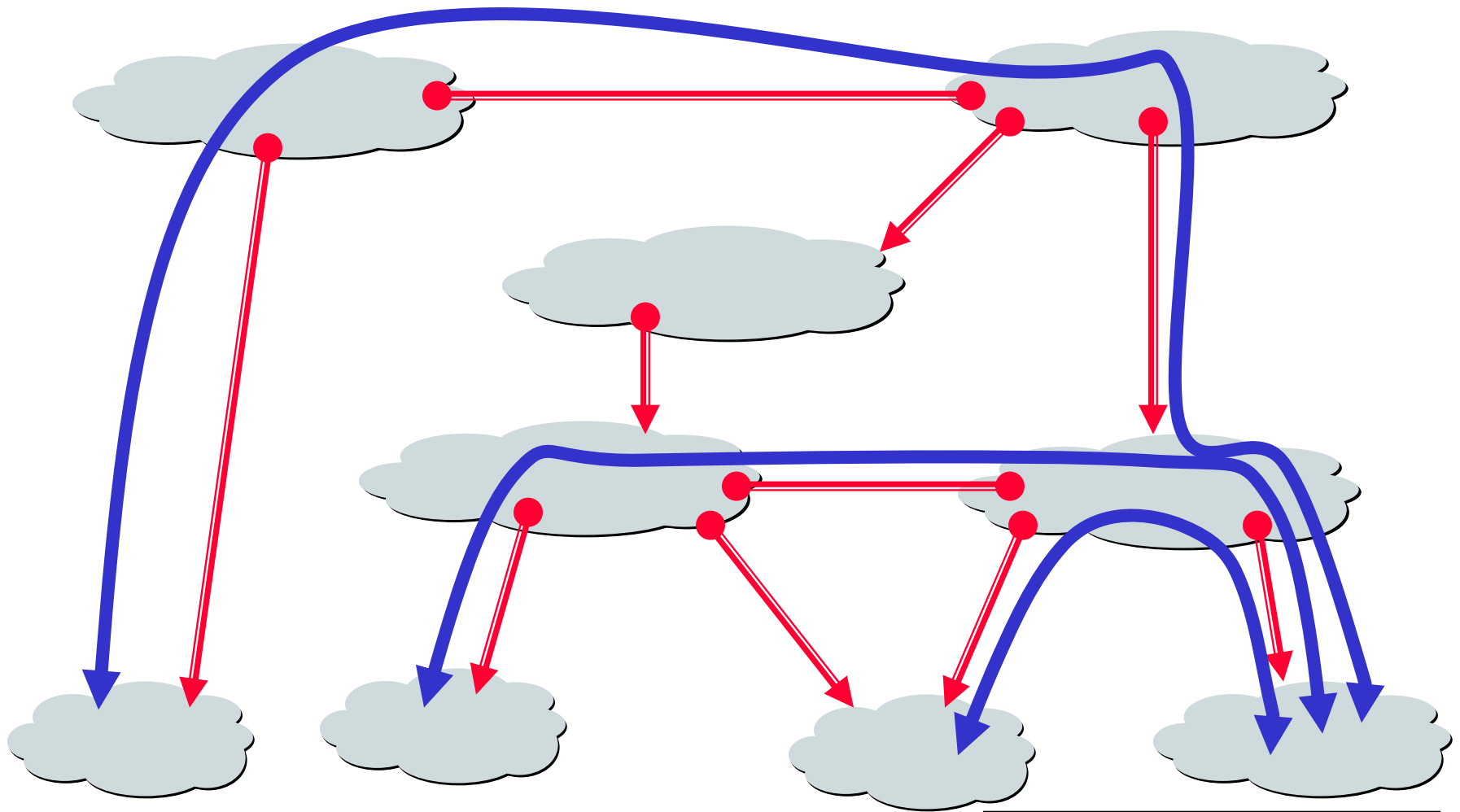


traffic allowed

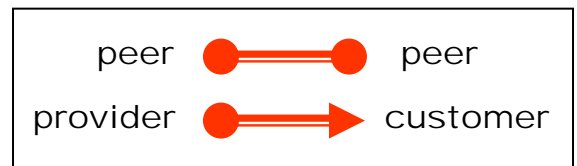


traffic NOT allowed

# Peering Provides Shortcuts



Peering also allows connectivity between the customers of "Tier 1" providers.



# BGP: Path Vector Protocol

- Distance vector algorithm with extra information
  - For each route, store the complete path (ASs)
  - No extra computation, just extra storage
- Advantages:
  - can make policy choices based on set of ASs in path
  - can easily avoid loops

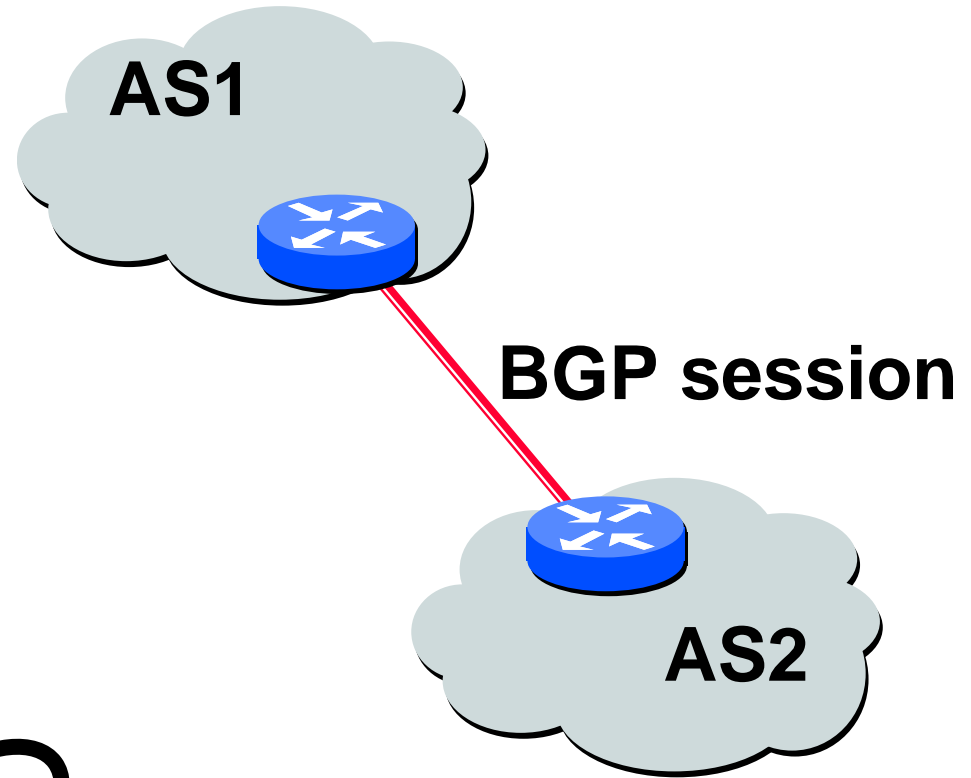


# BGP Operations (Simplified)

Establish session on  
TCP port 179

Exchange all  
active routes

Exchange incremental  
updates



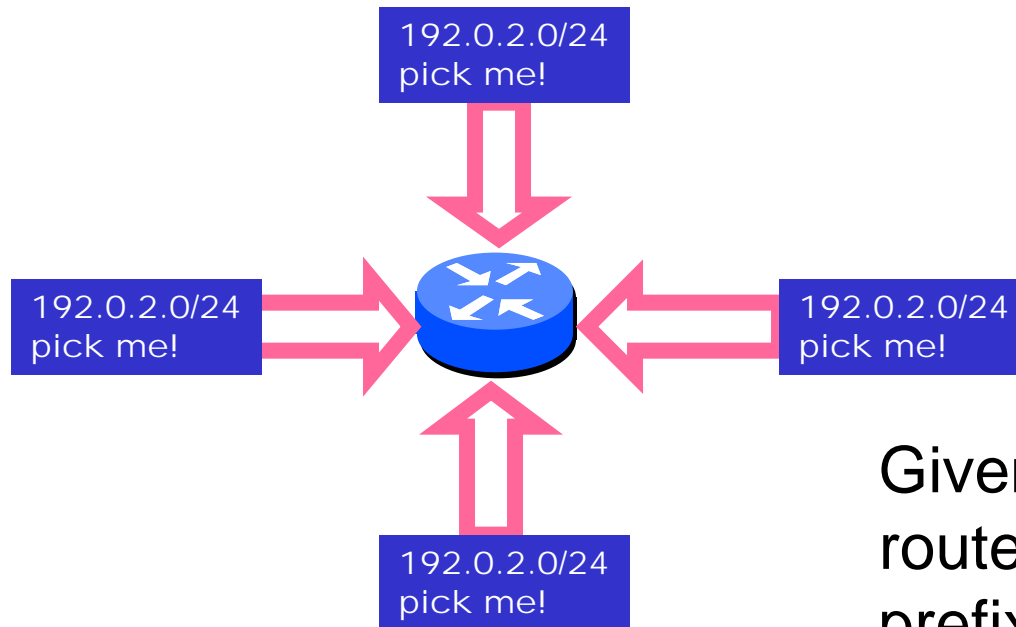
While connection  
is ALIVE exchange  
route UPDATE messages

# Four Types of BGP Messages

- **Open** : Establish a peering session.
- **Keep Alive** : Handshake at regular intervals.
- **Notification** : Shuts down a peering session.
- **Update** : Announcing new routes or withdrawing previously announced routes.

Announcement  
=  
prefix + attributes values

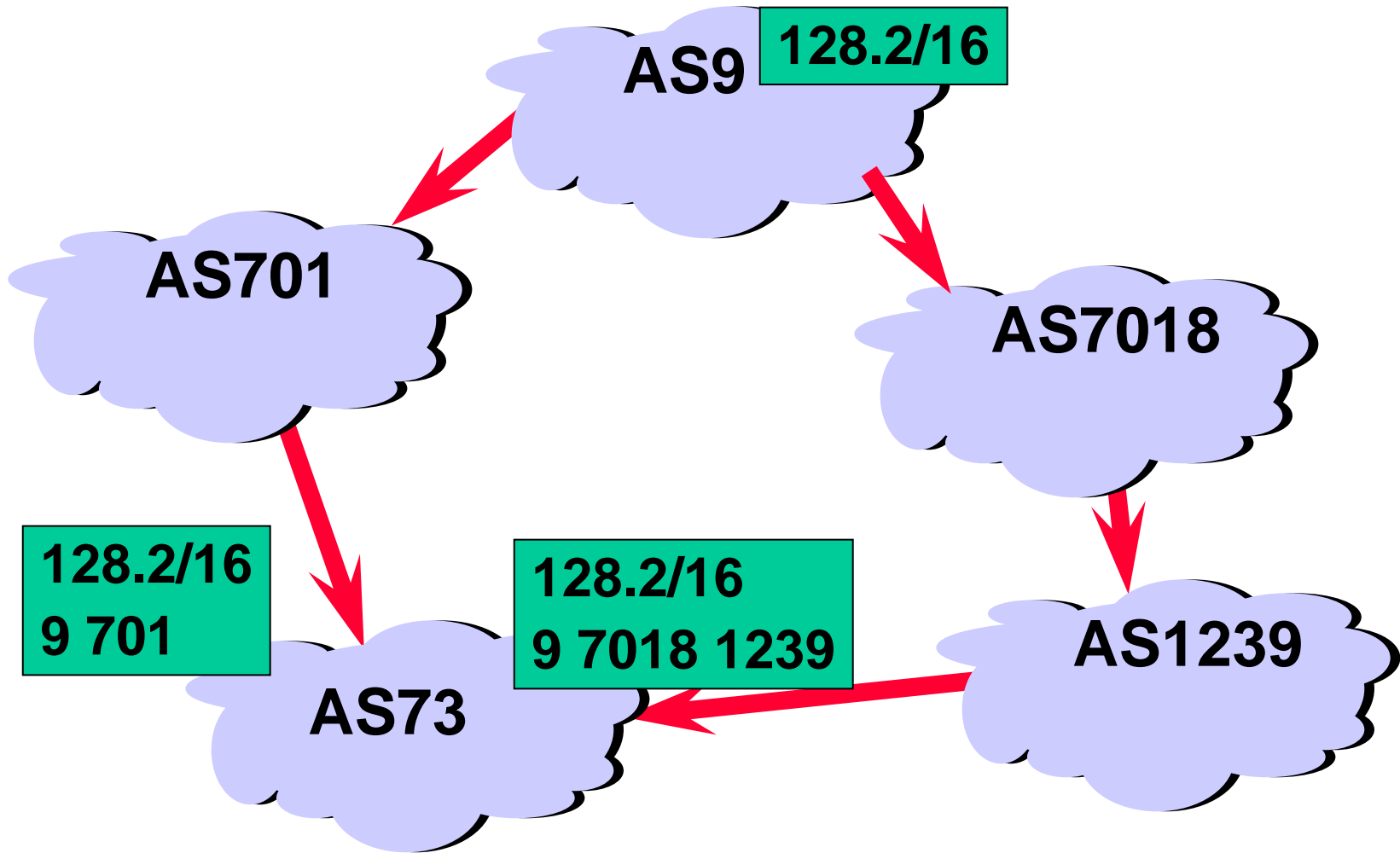
# Attributes are Used to Select Best Routes



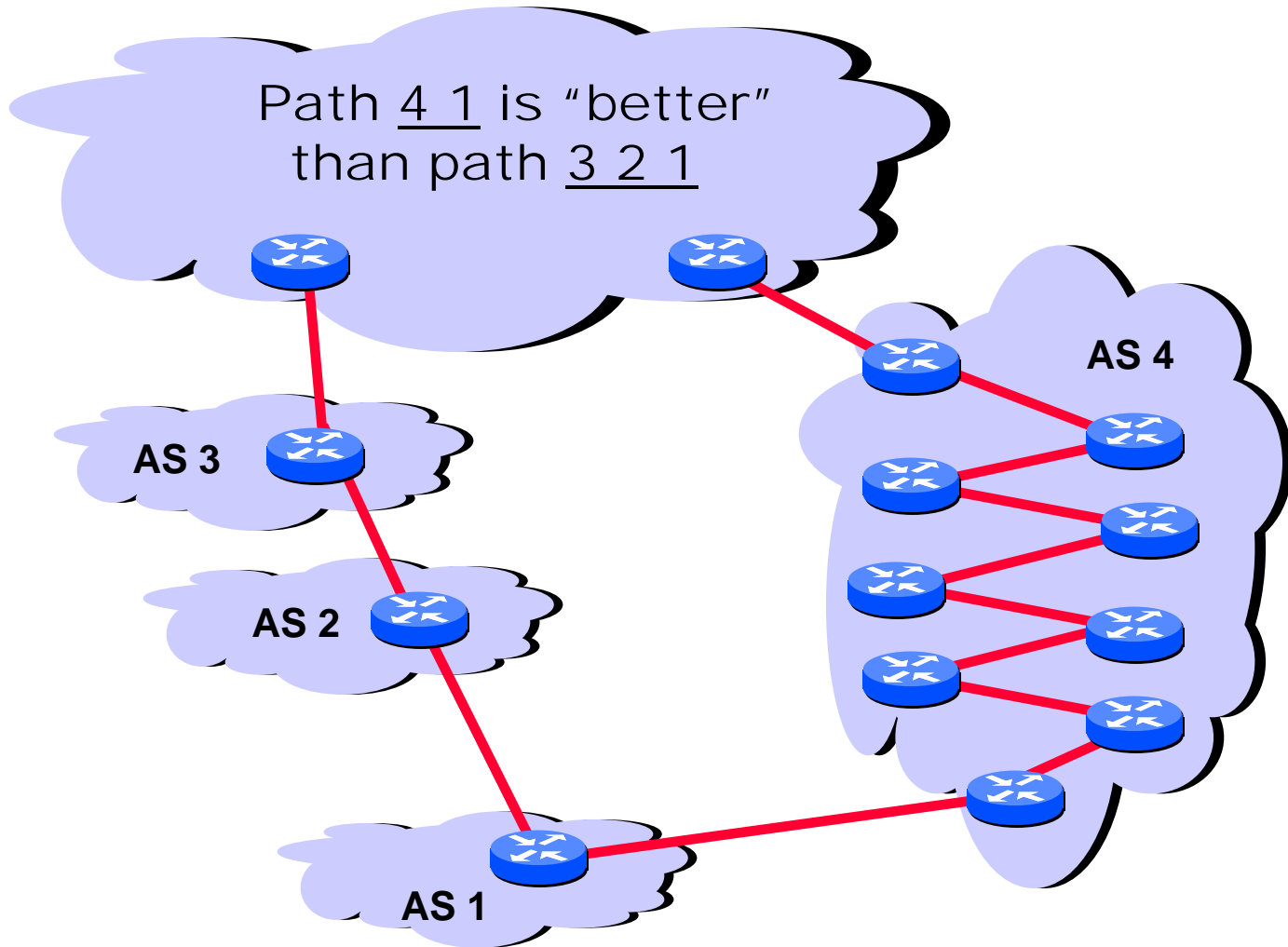
Given multiple routes to the same prefix, a BGP speaker must pick at most one best route

(Note: it could reject them all!)

# Example: Multiple AS Paths



# Shorter Doesn't Always Mean Shorter



## Implementing Customer/Provider and Peer/Peer relationships

- What you announce determines what route can be used by whom
- Enforce transit relationships
  - Outbound route filtering
- Enforce order of route preference
  - provider < peer < customer

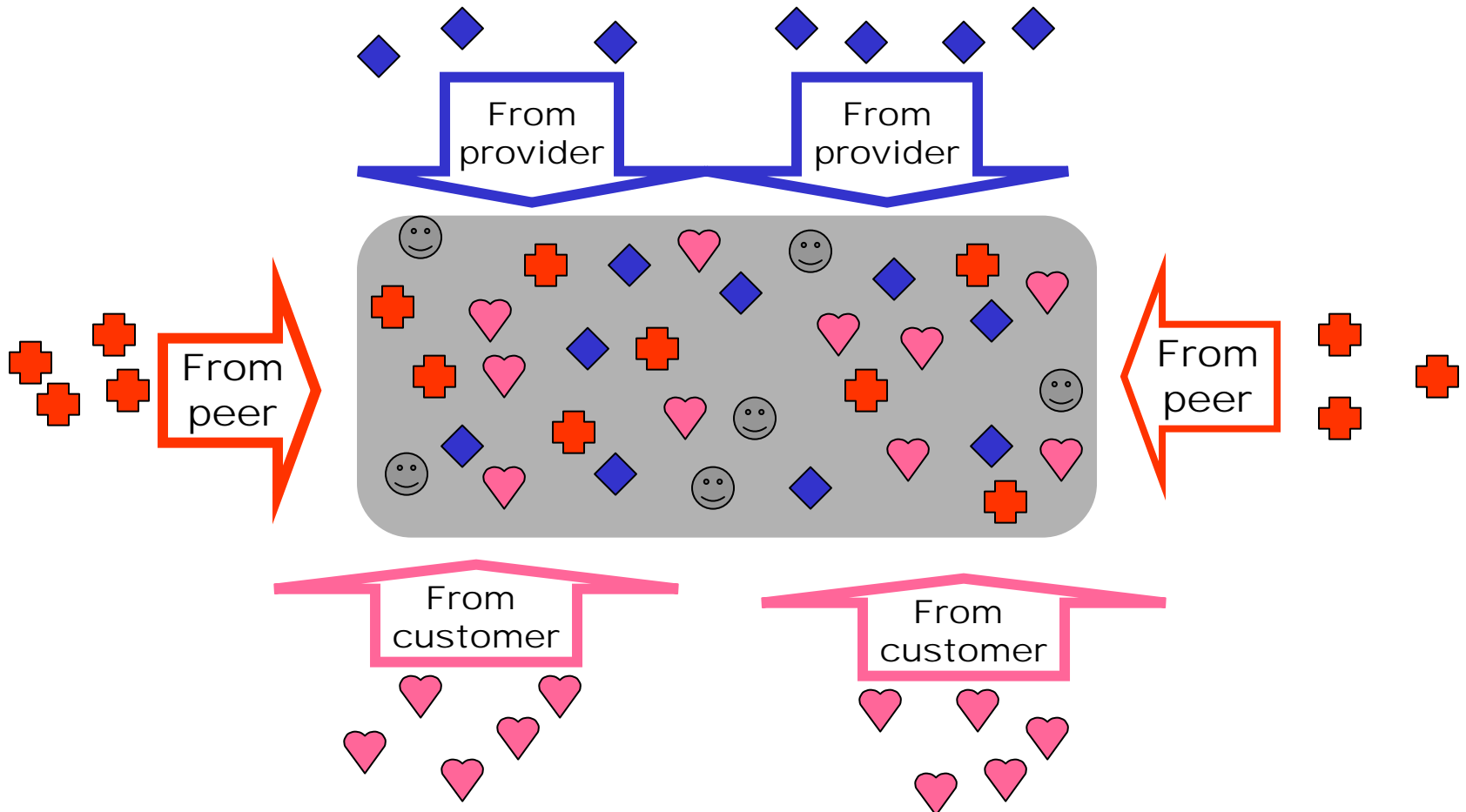
# Import Routes

◆ provider route

⊕ peer route

♥ customer route

☺ ISP route



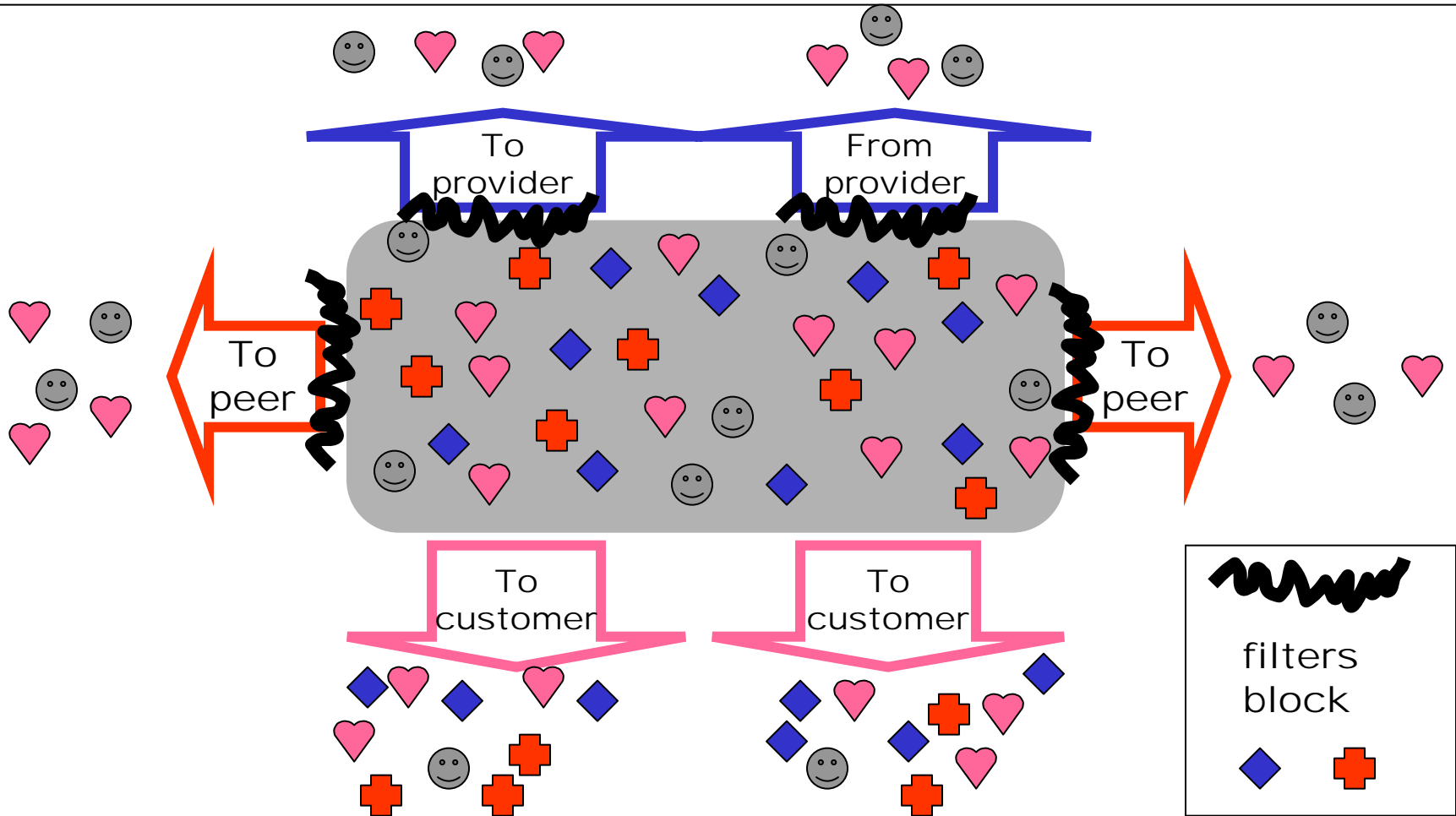
# Export Routes

◆ provider route

⊕ peer route

♥ customer route

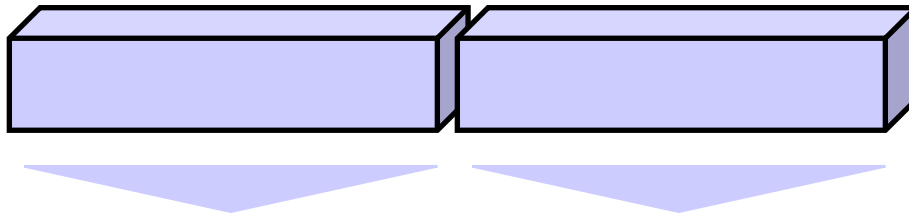
☺ ISP route





# How Can Routes be Colored? BGP Communities!

A community value is 32 bits



By convention,  
first 16 bits is  
ASN indicating  
who is giving it  
an interpretation

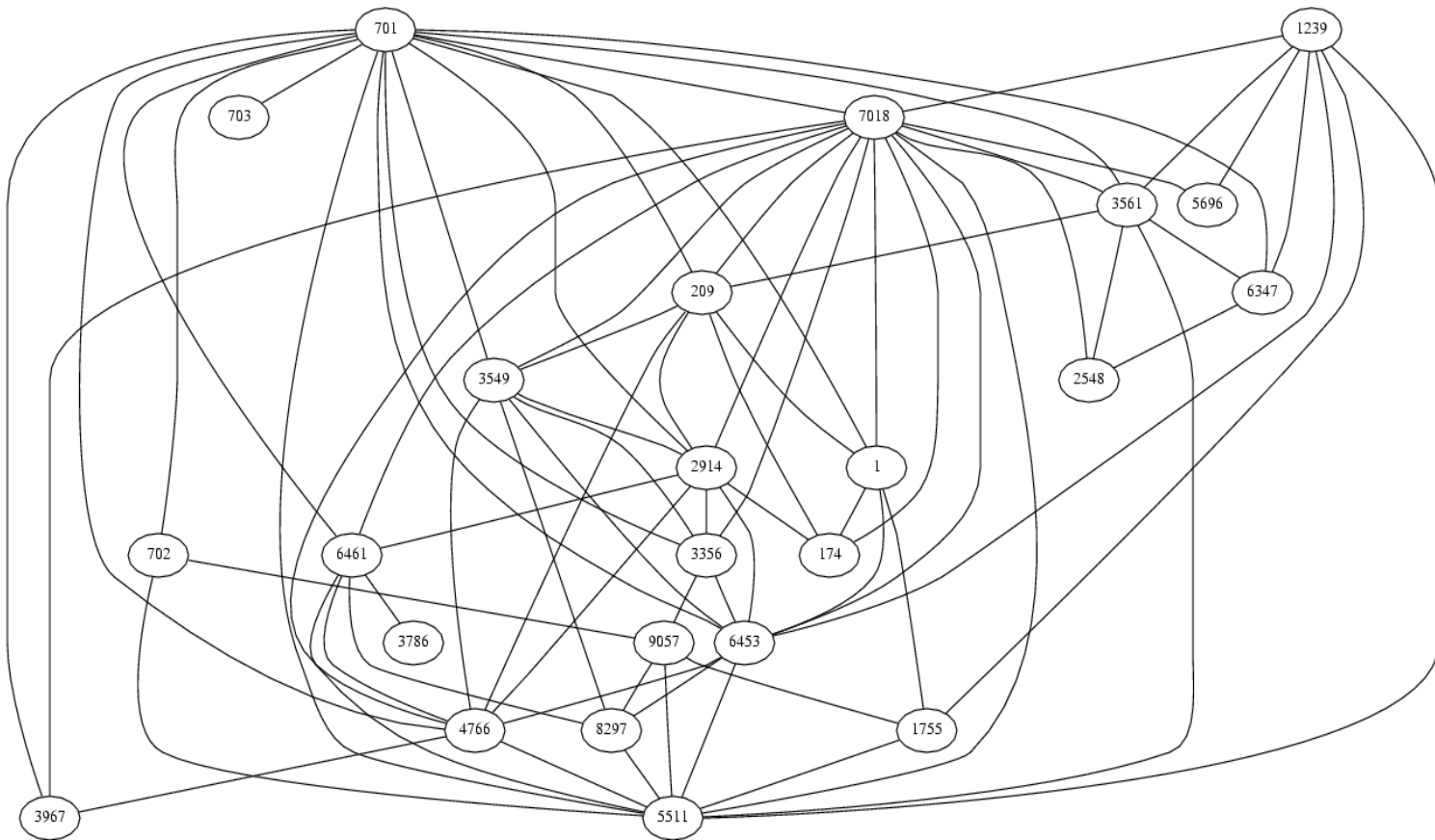
community  
number

Used for signaling  
within and between  
ASes

Very powerful  
**BECAUSE** it  
has no predefined  
meaning

**Community Attribute = a list of community values.  
(So one route can belong to multiple communities)**

# Example AS Graph



The subgraph showing all ASes that have more than 100 neighbors in full graph of 11,158 nodes. July 6, 2001. Point of view: AT&T route-server

Does not reflect true topology

# BGP Issues

- BGP designed for policy not performance
- Susceptible to router misconfiguration
  - Blackholes: announce a route you cannot reach
- Incompatible policies
  - Solutions to limit the set of allowable policies

# More Issues

- Scaling the I-BGP mesh
  - Confederations
  - Route Reflectors
- BGP Table Growth
  - 140K prefixes and growing
  - Address aggregation (CIDR)
  - Address allocation
- AS number allocation and use
- Dynamics of BGP
  - Inherent vs. accidental oscillation
  - Rate limiting and route flap dampening
  - Lots and lots of redundant info
  - Slow convergence time